

UE ELA

PO 5 78763

RANG-TR-78-185; Volume V (07 five) Final Technical Report October 1978

Marie 1 40 57 8 75

DAYESTAN SOFTWIRE PREDICTION MODELS Summery of Technical Progress

Anrie L. Hal

Syracuse University

Market Strate

Approved for public release; distribution unlimited.

Color and resident matter charges

0 04 20 m

This report has been reviewed by the RADC Information Office (DI) and is releasable to the National Technical Information Service (NTIS). At NTIS it will be releasable to the general public, including foreign mations.

RADC-TR-78-155, Volume V (of five) has been reviewed and is approved for publication.

Marian.

alan N. Sukest

Project Engineer

APPROVED:

Mare Brume

WENDALL C. BAUMAN, Colonel, USAF Chief, Information Sciences Division

FOR THE COMMANDER: JAN G. Kurs

JOHN P. HUSS

WAR ON THE RESIDENCE OF STREET

to the first of the second of

Acting Chief, Plans Office

If your address has changed or if you wish to be removed from the RADC mailing list, or if the addresses is no longer employed by your organization, please notify RADC (ISIS) Griffies AFS BY 13441. This will assist us in maintaining a current mailing list.

Do not return this copy. Retain or destroy.

UNCLASSIFIED SECURITY CLASSIFICATION OF THIS PAGE (When Date Entered) READ INSTRUCTIONS BEFORE COMPLETING FORM PEPOPT DOCUMENTATION PAGE RADC TR-78-155 final Technical Report. Volume X BAYESIAN SOFTWARE PREDICTION MODELS . Jan 76- Apr 784 Summary of Technical Progresse Technical Report No 78-54 AUTHOR(a) F30602-76-C-0097 Amrit L. Goel PERFORMING ORGANIZATION NAME AND ADDRESS Syracuse University Syracuse NY 13210 55811 11. CONTROLLING OFFICE NAME AND ADDRESS Rome Air Development Center (ISIS) Griffiss AFB NY 13441 94 4. MONITORING AGENCY NAME & ADDRESS(II different from Controlling Office) 15. SECURITY CLASS. (of this report) Same UNCLASSIFIED 15a. DECLASSIFICATION/DOWNGRADING 16. DISTRIBUTION STATEMENT (of this Report) Approved for public release; distribution unlimited. 17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report) Same 18. SUPPLEMENTARY NOTES RADC Project Engineer: Alan N. Sukert (ISIS) 19. KEY WORDS (Continue on reverse side if necessary and identify by block number) Software Error Prediction Software Demonstration Testing Imperfect Debugging Software Error Correction Imperfect Maintenance Bayesian Models Software Modeling Bayesian Inference

DD 1 JAN 73 1473

test plans is described.

339600

This report summarizes the technical activities pursued under Contract

F30602-76-C-0097, Bayesian Software Prediction Models, with Rome Air Development Center. Research work, discussed in previous volumes, discussing imperfect debugging and imperfect maintenance software performance models, is summarized and some additional work in development of software reliability demonstration

ABSTRACT (Continue on reverse side if necessary and identify by block number)

UNCLASSIFIED

SECURITY CLASSIFICATION OF THIS PAGE (When Date Entered)

TABLE OF CONTENTS

NTIS		Section	*
BOC	Buff	Section	
JUSTIFICATION	_		
BY	Van Ade	ITY MODE	
DISTRIBUTION/A	AHLKOL	מנו שוו ששני	1
		or SP:C	

1.	INTR	DUCTION	2
2.	A SO	TWARE PERFORMANCE MODEL UNDER IMPERFECT DEBUGGING .	4
26	2.1	Model and Main Results	4
		2.1.1 Distribution to time to a specified number of remaining errors	5
		2.1.2 Probability distribution of a given number of remaining errors at time t	6
		2.1.3 Expected number of total and imperfect debugging errors	6
		2.1.4 Reliability function	7
		2.1.5 Gamma approximation	7
		2.1.6 Numerical examples	8
	2.2	Analysis of Total and Imperfect Debugging Errors in a Real Time Control System	11
3.		ABILITY ANALYSIS OF SOFTWARE SYSTEMS UNDER	12
	3.1		12
	3.2	Numerical Example	14
	3.3	A Nomogram for the Expected Time to a Specified Number of Errors and to Determine Manpower	
		Requirements	14
4.	- 10-21 II-10 TO TO TO THE	IAN AND CLASSICAL INFERENCE FOR THE IMPERFECT GING AND MAINTENANCE MODELS	20
	4.1	Maximum Likelihood Method	20
	4.2	Bayesian Inference	22

5.	BAYESIAN	SOFT	WARE	CORR	ECTION	LIMI	r POL	ICIES	s	•		4.	•	24
6.	SOFTWARE	RELI	ABIL	TY D	EMONST	RATIO	N TES	T PL	ANS.		•			26
7.	COMPUTER	PROG	RAMS.	• •		•	ų • į	• •		٠			٠	27
SEL	ECTED REI	FERENC	ES.							•			•	28
APP	ENDIX A	AN AN REAL-	NALYS	IS OF	RECUE	RENT	SOFT	VARE	ERRO	RS •	IN	. A		A-1
APP	ENDIX B	SOFT	VARE	RELIA	BILITY	DEMO	NSTRA	TION	TES	т.	.0		٠	B-1
APP	ENDIX C	COMP	JTER :	PROGE	RAMS	•	98	13,542		•	•		٠	C-1
		c.1	Prog Mode	rams 1 (Se	for thection	ne Imp	erfe	ct De	bugg	in	g •		•	c-1
		C.2	Prog. Debu	rams gging	for Si	mulat	ion o	of Im	perf	ect	t		•	C-1
		C.3	Prog Mode	rams 1 of	for the	ne Imp	erfe	et Ma	inte	naı	nce •		74	C-1
		C.4	Prog Limi	ram f	or Bay	esian (Sect	Soft	ware	Cor	red	cti	on .	16	C-2

1700 1700 2700 1000 1703 200

PANHODINAL P

TANTHEMEN SHE THE TOWNSHIP LANTERALD THE THE MADRITURE WAS MADRITURED BY THE STATE OF THE STATE

LIST OF FIGURES

a many 30 1214

Figure	the first of the second	Page
2.1	Probability Distribution of Time to n ₀ Remaining Errors	9
2.2	Expected Number of Remaining Errors versus Time	10
3.1	Plots of State Occupancy Probabilities and Software System Availability	15
3.2	Expected Number of Software Errors Detected and Corrected by Time t	16
3.3	A Nomogram for the Expected Time to a Specified Number of Errors	18
5.1	Sequence of Corrective Actions in Operational Phase	25
A.1	Joint confidence regions for N and λ for p= \hat{p}	A-5
A.2	Actual and fitted SPRs by month	A-6
A.3	Plots of the actual and predicted number of remaining errors	A-7
B.1	Contours of $\bar{\alpha}$, β^* for Design of Test Plans	B-18
B.2	Contours of $\overline{\alpha}$, β^{**} for Design of Test Plans	B-19

100 A

LIST OF TABLES

Ta	ble																					Page
	A.1	Error Data	by Mont	h	•	٠			•						•							A-2
	A.2	A Summary	of Error	. [at	ta	Ar	nal	Lys	ses	3.	u.i				•				d.		A-4
	c.1	Subroutine	FRSTPT		•	•	٠	•	•			•			•			•	•	•	•	C-2
	C.2	Subroutine	FPT2 .	•	•	1.5	•	•	•	•	•	•	٠	•				٠	•	•	•	C-3
	c.3	Subroutine	State.	•		E CA	.•				•	•		•	4.5	•						C-4
	C.4	Subroutine	Mean .	•	a•1	6. 3	•1	-0		•		•		•)	•	•	٠.		÷.,		na.	C-5
	C.5	Subroutine	TBF	•		•					•	٠	•	٠	•	٠	٠	٠			•	C-6
	C.6	Subroutine	MDGAMM	•							•		•	•	•	٠	•	•		•		C-7
	C.7	Subroutine	GAMMA.		•	•	9.0	•		•			•	•	•	•	•	64.2	4			C-9
	C.8	Subroutine	SMLT .	•		•	•	•			•	•		.0	c.		•				ST.	c-11
	c.9	Subroutine	MLE	•	•			•1	•	٠						•	•	•	1.8	•	•	C-12
	C.10	Subroutine	BAYES.			•		•	•	•	•		•	٠		•	•	•	•	•	•	C-14
	c.11	Subroutine	GGUB .		•				٠		٠		•	•		٠	•		•		•	C-16
5	C.12	Subroutine	COMP .	•	٠		•	•	:	•		•	•	•	•	•	•	•		•	•	C-18
	c.13	Subroutine	FIRST.	•	•	•	•	٠	•	•	•	•	•		•			•	•		•	C-19
	C.14	Subroutine	STT	•				•	•					•	•		•	•			•	C-20
	C.15	Subroutine	AVAIL.	•	•			•	•	•	•	•							•			C-21
	C.16	Subroutine	EXPCT.		•				•	•			•	•	•			•	•		•	C-22
	C.17	Subroutine	MDGAMM			•	•			•	٠	•										C-23
	C.18	Subroutine	GAMMA.	•	•		•	•							•						•	C-25
	C.19	Subroutine	MDL1 .	•			•											•		•	•	C-27
	C.20	Subroutine	MDL2 .	•	•					•		•	•	•	•	•	•			•	•	C-28
	C.21	Subroutine	DATA1.	•	•	٠	•				•		•		•	•	•		•	•	٠	C-30
	C.22	Subroutine	GGUB .	•	•						•		•					•	•		•	C-31
	C.23	Subroutine	OPTMM.																			C-32

EVALUATION

The necessity for more complex software systems in such areas as command and control and intelligence has led to the desire for better methods for predicting software errors and reliability to insure that software produced is of higher quality and of lower cost. This desire has been expressed in numerous industry and Government sponsored conferences, as well as in documents such as the Joint Commanders' Software Reliability Working Group Report (November 1975). As a result, numerous efforts have been initiated to develop and validate mathematical models for predicting such quantities as the number of remaining errors in a software package and the time to achieve a desired level of reliability. In addition, efforts have been initiated to develop better methods for determining when a software package should be released to a potential user. However, these efforts have not produced measures with the desired accuracy or confidence for general applicability.

This effort was initiated in response to this need for developing better and more accurate software error prediction and demonstration tests and fits into the goals of RADC TPO No. 5, Software Cost Reduction in the subthrust of Software Quality (Software Modeling). This report summarizes the development of mathematical models for predicting quantities such as the expected number of errors during both software development and software maintenance. These models assume errors are not corrected with probability 1, i.e. imperfect development and maintenance. The report also describes the development of statistical tests for determining whether a software package should be accepted or rejected after completion of testing. The importance of these developments is that they represent the first attempt to develop both software error prediction models that incorporate imperfect debugging and thus more closely reflect the actual software error detection and correction process, and software demonstration tests that allow better statistical criteria for accepting a software package.

The theory and equations developed under this effort will lead to much needed predictive measures for use by software managers in more accurately tracking software development projects in terms of stated error and reliability objectives. In addition, the associated confidence limits and other related statistical quantities developed under this effort will insure more widespread use of these techniques. The acceptance criteria developed will permit better control of the release of software packages so that software is not given to a potential user before it is ready for operational usage. Finally, the measures developed under this effort will be applicable to current software development projects and thus help to produce the high quality, low cost software needed for today's systems.

ALAN N. SUKERT Project Engineer

ABSTRACT

This report provides a summary of the technical activities pursued under Contract F30602-76-C-0097 with RADC during January 1976-April 1978. Research work discussed in previous reports under this contract is summarized and some additional work is described. Also included is a brief description of research in progress.

ill passible we eschalp the ere browned to the edgic last offe

to govern the manner, the armental part actor . The commendant manner

THE RESIDENCE OF STREET

1. INTRODUCTION

During the past ten years the field of software engineering has grown considerably in importance and scope. A primary motivation for this growth has come from an ever increasing cost of developing and maintaining software systems. This is specially true for the DOD which needs high quality, low cost software for its operations. As a result of this increased importance, various fields within software engineering are maturing into disciplines of study and research, for example, software design techniques, structured programming and other improved programming methodologies, program testing and debugging techniques, software performance modelling, and techniques of program validation and verification. The ultimate objective of studies in all these fields is to develop tools that will be useful in the design, development and operational phases of the software life cycle. The objective of studies dealing with software error analysis and modelling has been to develop analytical tools which can be used for improving software performance. Such studies can be classified into one (or both) of two categories. In the first category the emphasis is on the analysis of software error data collected from small or large projects, during development and/or operational phases. Studies in the second category are primarily aimed at the development of analytical models which are then used to obtain the reliability and other quantitative measures of software performance.

Typical of the first category are the studies by Akiyama [1]
Belady and Lehman [3], Fries [6], Endres [5], Baker [2], Motley et al
[18], Miyamoto [16], Willman et al [35], Schneidewind [26],

Shooman et al [29], Sukert [30,31], Rye et al [24], Thayer et al [32], and Wagoner [34]. These studies range in size from an analysis of small data sets (108 errors), e.g. Wagoner [34], to analysis of large sets (3500 errors), e.g. Thayer et al [32] and encompass data from an on-line system [16], an operating system [3], to that from the Apollo project [24].

In the second category of papers, several models have been proposed and studied during the last six years. These include 'exponential type' models of Shooman [28], Jelinski and Moranda [11,12], and Schick and Wolverton [25]; models based on the non-homogeneous Poisson process proposed by Goel and Okumoto [9] and Schneidewind [27], and a Bayesian model by Littlewood and Verrall [15]. Halstead [10] has developed a theory based on 'software physics' for various measures of the performance of a software system.

Musa [19] has introduced a model which is based on a large number of parameters derived from the software system being modelled.

Trivedi and Shooman [33] consider a Markov model in which they incorporate the time spent for removal of errors.

Most of the above studies assume that errors are removed with certainty when detected. The purpose of the investigation summarized in this report is to develop and study models for software performance which account for the probabilistic nature of the programmer's action during debugging and operational phases of the software system, to provide a methodology for classical and Bayesian inference for various quantitative measures of performance, to develop optimum Bayesian software correctional limit policies, and to develop software reliability demonstration test plans. Results of these studies are summarized in Sections 2 through 6.

2. A SOFTWARE PERFORMANCE MODEL UNDER IMPERFECT DEBUGGING

The purpose of this modelling effort was to establish quantitative measures for software systems by incorporating the probabilistic nature of the programmer's actions during the debugging phase. Such occurrences have been termed as recurrent errors by Fries [6], Thayer [32] and Willman et al [35], erroneous debugging by Miyamoto [16], bad fixes by Jones [13] and accounted as an error reduction factor by Musa [19]. In this study we call them the imperfect debugging errors.

Even though the presence of the imperfect debugging phenomenon has been known, no published model, with the possible exception of Musa's error reduction factor, provides an explicit way to account for it in software performance prediction. The model and other related results for this topic are summarized below. Details of this work are given in [7]. These results are useful for software development personnel in establishing manpower requirements to achieve a desired quality in the system before it is released for operational use. Trade-off studies between cost of debugging and software quality can be undertaken using the results given below.

2.1 Model and Main Results

The following parameters are used for model development and analysis:

- N = the initial number of errors in the software system at the beginning of the debugging activity
- p = probability that the error causing a software failure is removed when detected

q = 1-p, the probability of imperfect debugging λ = software error occurrence rate

Let a random variable X(t) denote the number of errors remaining in the system at time t. Then X(t) describes the state of the system at time t. We consider the stochastic process $\{X(t), t \ge 0\}$ to be a semi-Markov process with the one-step transition probability, $Q_{ij}(t)$, the probability that the next failure resulting in j remaining errors will be by time t when a software system has i remaining errors at time zero. Then

$$Q_{ij}(t) = \begin{cases} p(1-e^{-i\lambda t}) & \text{if } j=i-1 \\ q(1-e^{-i\lambda t}) & \text{if } j=i \end{cases}$$

If we start with N errors at t=0, we are interested in the expressions for various quantities that describe the software system performance. These quantities are given below.

2.1.1 Distribution to time to a specified number of remaining errors

Let $G_{N,n_0}^{(t) \equiv P} (T_{N,n_0}^{\leq t})$ be the cdf of the time $T_{N,n_0}^{(t)}$ required to obtain a software system with n_0 remaining errors, $n_0=0,1,2,\ldots,N-1$. Then

$$G_{N,n_0}(t) = \sum_{j=1}^{N-n_0} B_{N,j,n_0}\{1-e^{-(n_0+j)p\lambda t}\},$$

where

$$B_{N,j,n_0} = \frac{N!}{n_0!j!(N-n_0-j)!} (-1)^{j-1} \cdot \frac{j}{n_0+j}.$$

The expected time required to obtain a software system with $\mathbf{n}_{\mathbf{0}}$ remaining errors is given by

$$E[T_{N,n_0}] = \sum_{j=1}^{N-n_0} B_{N,j;n_0}/(n_0+j)p\lambda$$
.

2.1.2 Probability distribution of a given number of remaining errors at time t

Probability that there are n_0 remaining errors at time t is

$$P_{N,n_0}^{(t)}(t) = P\{X(t)=n_0 | X(0)=N\}$$

= $G_{N,n_0}^{(t)-G_{N,n_0}-1}(t)$, $n_0=0,1,2,...,N$,

where

$$G_{N,N}(t)=1$$

and

$$G_{N,-1}(t) = 0$$
.

Also, the expected number of errors remaining at time t

$$E[X(t)|X(0)=N] = Ne^{-p\lambda t}.$$

2.1.3 Expected number of total and imperfect debugging errors

The expected total number of errors, $M_N(t)$, and errors due to imperfect debugging, $D_N(t)$, during a debugging time period t are given by

$$M_N(t) = \frac{N}{p} (1-e^{-\lambda pt})$$
,

and

$$D_{N}(t) = q \cdot M_{N}(t) .$$

2.1.4 Reliability function

The software system reliability between (k-1)st and kth failures is given by

$$R_{k}(x) = \sum_{j=0}^{k-1} {k-1 \choose j} p^{k-j-1} q^{j} e^{-\{N-(k-j-1)\}\lambda x}.$$

2.1.5 Gamma approximation

The computation of the quantity B_{N,j,n_0} and hence G_{N,n_0} (t) becomes cumbersome when N is large. We have found that the following gamma approximation yields satisfactory results for large software systems with large values of N .

$$G_{N,n_0}(t) \approx 0^{\int \frac{1}{r(\alpha)} \cdot e^{-\beta x} \alpha x}$$
,

where the scale parameter $\,\beta\,$ and the shape parameter $\,\alpha\,$ are estimated as

$$\beta = p\lambda \frac{\sum_{j=n_0+1}^{N} \frac{1/j}{n}}{\sum_{j=n_0+1}^{N} \frac{1/j^2}{n}},$$

and

$$\alpha = \frac{\begin{cases} \sum_{j=n_0+1}^{N} 1/j \end{cases}^2}{\sum_{j=n_0+1}^{N} 1/j^2}.$$

2.1.6 Numerical examples

To illustrate the usefulness of the above results consider a software system with N=100, λ =0.02. The probability distributions of times to n_0 =0(1)10 remaining errors are obtained as above and are shown in Figure 2.1 for p=0.9. We see that at t=200, say, the probability of having zero errors in the system is approximately 0.1, of one error about 0.15, of 2 about 0.23 and so on. Plots of expected number of remaining errors at various times are given in Figure 2.2 for p=0.8(0.05)1.00. For p=0.9 and t=100, there will be about 18 errors left in the system. From a study of these plots, one can plan the resource requirements and also conduct trade-off studies between available resources and resulting product.

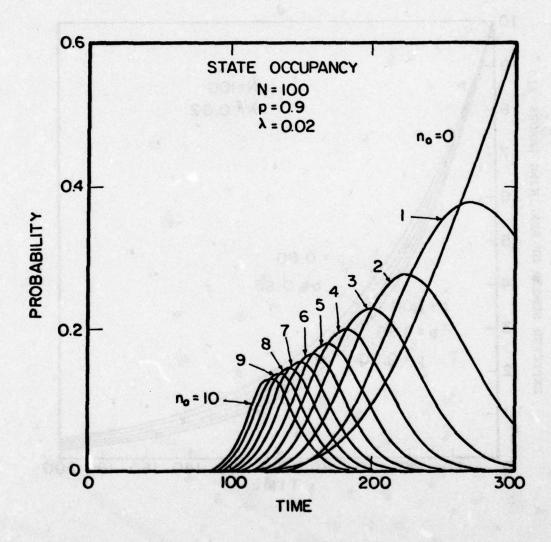


Figure 2.1 Probability Distribution of Time to n_0 Remaining Errors

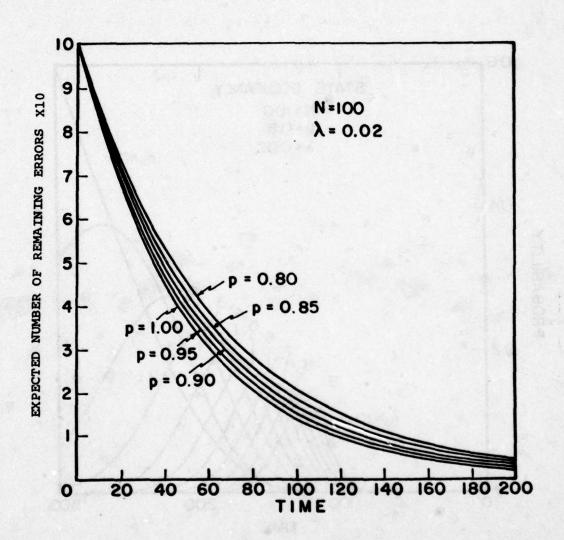


Figure 2.2 Expected Number of Remaining Errors versus Time

2.2 Analysis of Total and Imperfect Debugging Errors in a Real-Time Control System

The results described in Section 2.1 were used to analyze the software error data from a large-scale software project - a real-time control system for a land-based radar system developed by Raytheon Co. in a modular fashion, see Willman [35]. The data base was extracted from 2165 Software Problem Reports written against 109 operational software modules over the development phase. Details of this analysis are given in Appendix A.

3. AVAILABILITY ANALYSIS OF SOFTWARE SYSTEMS UNDER IMPERFECT MAINTENANCE

In this section we describe a model developed for the operational phase of the software system subject to imperfect error maintenance and also incorporate the time spent for error maintenance.

3.1 Model and Performance Measures

The following parameters are used in this model:

- N = Initial number of errors in the software system at the beginning of software operation
- p = Probability that the error causing a software failure
 is removed/maintained when detected.
- q = 1-p , the Probability of imperfect maintenance/removal
- λ_i = Software error occurrence rate per unit time when there are i errors in the system
- μ_i = Software error correction/maintenance rate per unit time when i errors remain in the system

Consider a stochastic process $\{X(t), t\geq 0\}$, whose states are defined as

$$X(t) = \begin{cases} i & \text{if the software system is operational while there} \\ & \text{are i errors in the system } (i=0,1,2,...,N) \\ & \text{D if the software system is down for error removal} \\ & \text{i.e., for maintenance.} \end{cases}$$

Then, the process $\{X(t), t\geq 0\}$ forms a semi-Markov process with the one step transition probability (the probability that the next up-down cycle, resulting in j remaining errors, will be completed

by time t when a software package has i remaining errors at time zero) given by

$$Q_{ij}^{(D)}(t) = \begin{cases} p(1-e^{-\lambda_{i}t}) \cdot (1-e^{-\mu_{i}t}) & \text{if } j=i-1 \\ q(1-e^{-\lambda_{i}t}) \cdot (1-e^{-\mu_{i}t}) & \text{if } j=i \\ 0 & \text{otherwise} \end{cases}$$

Based upon the above model, expressions for the following performance measures of the software system are derived by Okumoto and Goel [21]:

- (i) Distribution of time from N to a specified number of remaining errors n_0
- (ii) Expected time required for the system to go from N to n_0 errors, $E[T_{N,n_0}]$
- (iii) Probability of the system being operational at some time with n_0 remaining errors, $P_{N,n_0}(t)$.
- (iv) Software system availability, i.e., the probability of the system being operational at time t , $A(t) = \sum_{n_0=0}^{N} P_{N,n_0}(t)$
- (v) Probability that the number of errors remaining in the system is n , $P[\overline{N}(t)=n]$ and the expected number of errors at t , $E[\overline{N}(t)]$
- (vi) Expected number of errors detected and corrected by t , denoted by $\mathbf{M}_N^{\ D}(t)$ and $\mathbf{M}_N^{\ C}(t)$, respectively.

The above quantities are useful for software managers for estimating the time and manpower requirements to achieve a desired level of performance during the operational phase of the software system.

3.2 Numerical Example

For illustration purposes consider the case when $\lambda_1=i\lambda$, $\mu_1=i\mu$, N=100, p=0.9, λ =0.02 and μ =0.05. Using the expressions for various performance measures from [21], the plots of state occupancy probabilities and availability are given in Figure 3.1 and the plots of $M_N^D(t)$ and $M_N^C(t)$ are given in Figure 3.2. From a study of such plots for various values of λ , μ and μ , one can obtain adequate information about the behavior of the software system as well as about the resource requirements to achieve a desired level of performance. Thus, if μ is known to be 0.9, and the values of μ and μ that can be provided for, then the system availability at t=300 will be approximately 0.5. If this is not satisfactory, one has to provide additional or better resources that will yield better values for one or more parameters.

3.3 A Nomogram for the Expected Time to a Specified Number of Errors and to Determine Manpower Requirements

We present a simple nomogram to calculate the expected time required to remove a specified number of errors from a software system which will satisfy the desired performance requirements. We consider the case when $\lambda_i^{=i\lambda}$ and $\mu_i^{=i\mu}$, i.e., when the error detection and error correction rates are proportional to the number of remaining errors with λ and μ , respectively, the constants of proportionality. Letting the ratio $\lambda/\mu=\rho$, we have

$$E[T_{N,n_0}] = \frac{1+\varrho}{\rho_{\lambda}} \sum_{i=n_0+1}^{N} \frac{1}{i}$$
.

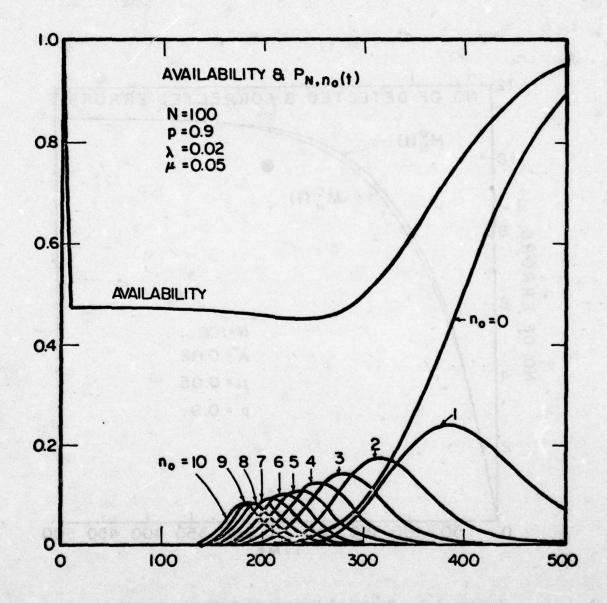


Figure 3.1 Plots of State Occupancy Probabilities and Software System Availability

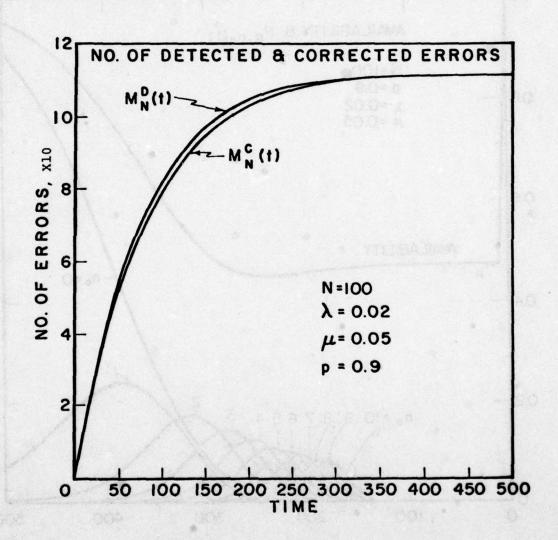


Figure 3.2 Expected Number of Software Errors
Detected and Corrected by Time t

For a large software system, this can be approximated as

$$E[T_{N,n_0}] \approx \frac{(1+\rho)}{p\lambda} \log(\frac{N+1}{n_0+1})$$
.

Letting $\phi(\phi,\rho) = (1+\rho)\phi$, where $\phi = \log(\frac{N+1}{n_0+1})$, we get

$$E[T_{N,n_0}] = \frac{\phi(\phi,\rho)}{p\lambda}.$$

A nomogram which gives the contours of $\phi(\phi,\rho)$ in the (ϕ,ρ) phase is given in Figure 3.3. To illustrate the use of this nomogram, consider the case when N=100, p=0.9, λ =0.02 per day, μ =0.05 per day and the desired value of n_0 is 10. We proceed as follows.

Step 1. Compute
$$\rho = \frac{\lambda}{\mu} = \frac{0.02}{0.05} = 0.4$$
 and $\phi = \log(\frac{N+1}{n_0+1}) = \log(\frac{100+1}{10+1}) = 0.963$.

Step 2. Corresponding to ϕ =.963 and ρ =0.4, from Figure 3.3 the value of ϕ (.963,0.4) is 3.2.

Step 3. Compute the value
$$E[T_{100,10}] = \frac{\phi(\phi,\rho)}{p\lambda} = \frac{3.2}{(0.9)(0.02)} = 177.8$$

Thus, for the given conditions the expected time required to go from 100 to 10 errors is 177.8 days.

Now we consider another application of this nomogram to show how it can be used in determining the manpower requirements for a specified objective of remaining errors. Suppose we want to go from N=1000 to n_0 =10 errors in $E[T_{1000,10}]$ =500 days when the error occurrence rate is λ =0.01 errors per day and the probability p of perfect correction is 0.95. We are interested in determining the manpower requirement to accomplish this objective.

The first thing to determine is the value of $\,\mu\,$ that will satisfy these requirements. We know that

$\Phi(\phi, P) = 2.5(0.5)25$

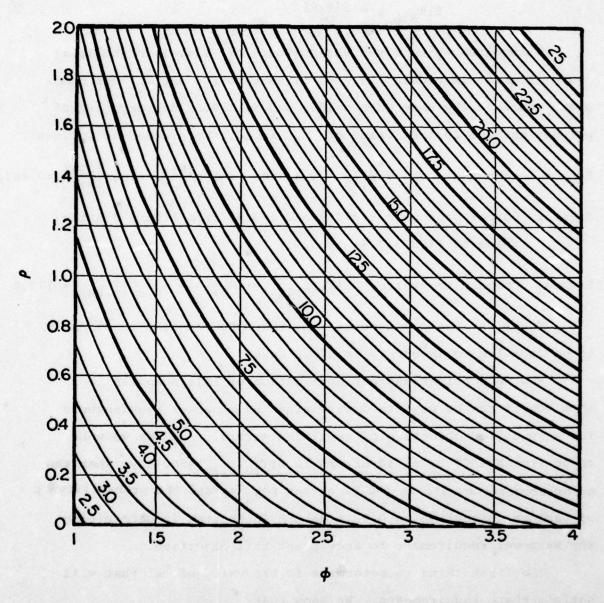


Figure 3.3 A Nomogram for the Expected Time to a Specified Number of Errors

$$\phi(\phi,\rho) = p\lambda E[T_{N,n_0}]$$

and hence

$$\phi(\phi,\rho) = (.95)(.01)(750) = 7.125.$$

From the nomogram in Figure 3.3, for $\Phi(\phi,\rho)=7.125$ and $\phi=\log(\frac{1001}{11})=1.96$, the value of $\rho=0.58$ and hence $\mu=\lambda/\rho=0.0172$ errors/day. If the error removal rate per person per day is 0.01, then we will need 1.7 people to meet the desired objective.

4. BAYESIAN AND CLASSICAL INFERENCE FOR THE IMPERFECT DEBUGGING AND MAINTENANCE MODELS

In this section we describe two methods for statistical inference of the parameters of the models described in Sections 2 and 3. The first one is the classical approach based on maximum likelihood estimation and the second is a Bayesian approach based on the prior distributions of the unknown parameters. The parameters under consideration are the initial number of software errors N, the error occurrence rate for each error λ , and the probability of perfect debugging p. An additional parameter for the imperfect maintenance model is μ . We give only the method for the model of Section 2. The same procedure can be used for the model of Section 3.

The available data for estimation purposes is generally given as $\underline{t}=(t_1,t_2,\ldots,t_n)$, the times between software failures and $\underline{y}=(y_1,y_2,\ldots,y_n)$, an indicator variable for imperfect debugging such that $y_i=1$ if the ith failure is caused by an error due to imperfect debugging and $y_i=0$ if the error is not due to imperfect debugging.

The maximum likelihood and Bayesian approaches to the estimation of the above parameters are summarized below. The details of the procedure are given in [20].

4.1 Maximum Likelihood Method

The likelihood function for N, p and λ is

$$L(N,p,\lambda | \underline{t},\underline{y}) = \prod_{i=1}^{n} \{N-p(i-1)\}e^{-\{N-p(i-1)\}\lambda t_{\underline{i}}} \cdot \{\frac{q(i-1)}{N-p(i-1)}\}^{\underline{y}_{\underline{i}}} \cdot \{\frac{N-(i-1)}{N-p(i-1)}\}^{1-\underline{y}_{\underline{i}}}$$

The maximum likelihood estimates $(\hat{N}, \hat{p} \text{ and } \hat{\lambda})$ are obtained as solutions of the simultaneous non-linear equations

$$\lambda \sum_{i=1}^{n} t_{i} = \sum_{i=1}^{n} \frac{1-y_{i}}{N-(i-1)}$$

$$\lambda \sum_{i=1}^{n} (i-1)t_{i} = \sum_{i=1}^{n} y_{i}/q$$

$$n/\lambda = \sum_{i=1}^{n} \{N-p(i-1)\}t_{i}.$$

The joint $100(1-\alpha)$ % confidence regions for N, p and λ are obtained from

$$\ell(\hat{N},\hat{p},\hat{\lambda}|\underline{t},\underline{y}) - \ell(N,p,\lambda|\underline{t},\underline{y}) = \frac{1}{2} \chi_{3,\alpha}^{2},$$

where

$$\ell(\cdot|\underline{t},\underline{y}) \equiv \log L(\cdot|\underline{t};\underline{y})$$
.

The estimated variance-covariance matrix for N, \hat{p} and $\hat{\lambda}$ is

$$\hat{\Sigma}_{cov} = \begin{bmatrix} r_{NN} & r_{Np} & r_{N\lambda} \\ r_{pN} & r_{pp} & r_{p\lambda} \\ r_{\lambda N} & r_{\lambda p} & r_{\lambda \lambda} \end{bmatrix}_{ \begin{subarray}{c} N=\hat{N} \\ p=\hat{p} \\ \lambda=\hat{\lambda} \end{subarray}}$$

where

$$r_{NN} = \sum_{i=1}^{n} 1/\{N-(i-1)\}\{N-p(i-1)\}$$

$$r_{Np} = r_{pN} = 0$$

$$r_{N\lambda} = r_{\lambda N} = \frac{1}{\lambda} \sum_{i=1}^{n} 1/\{N-p(i-1)\}$$

$$r_{pp} = \begin{cases} \frac{1}{q} \sum_{i=1}^{n} (i-1)/\{N-p(i-1)\} & \text{if } q \neq 0 \\ \infty & \text{if } q = 0 \end{cases}$$

$$r_{p\lambda} = r_{\lambda p} = -\frac{1}{\lambda} \sum_{i=1}^{n} (i-1)/\{N-p(i-1)\}$$

$$r_{\lambda \lambda} = n/\lambda^{2} .$$

4.2 Bayesian Inference

Now we describe a Bayesian approach for obtaining posterior point estimates and the highest posterior density (HPD) region for parameters N, p and λ .

The choice of the prior distribution for a parameter is governed by several factors. In our case we take the conjugate priors, which for N and λ are gamma distributions while for priors a beta distribution, i.e.,

$$P(N) = N^{\alpha-1} e^{-\beta N} , \qquad N>0$$

$$P(p) = p^{\alpha-1} (1-p)^{\rho-1} , \qquad 0 \le p \le 1$$

$$P(\lambda) = \lambda^{\mu-1} e^{-\gamma \lambda} , \qquad \lambda>0 .$$

By applying Bayes theorem the joint posterior distribution of N, p and λ for given priors and the data is obtained as

$$p(N,p,\lambda|t,y) \rightarrow p(N,p,\lambda)L(N,p,\lambda|t,y)$$
.

Let \hat{N} , \hat{p} and $\hat{\lambda}$ be the Bayesian point estimates for N, p and λ , respectively. That is, the point $(\hat{N},\hat{p},\hat{\lambda})$ is the mode of the joint posterior distribution $p(N,p,\lambda|\underline{t},\underline{y})$, i.e. it attains its maximum at $(\hat{N},\hat{p},\hat{\lambda})$. Then, \hat{N} , \hat{p} and $\hat{\lambda}$ are the values that satisfy

$$-\lambda \sum_{i=1}^{n} t_{i} + \sum_{i=1}^{n} \frac{1-y_{i}}{N-(i-1)} + \frac{\alpha-1}{N} - \beta = 0 ,$$

$$\lambda \sum_{i=1}^{n} (i-1) t_i - \sum y_i / (1-p) + \frac{\pi - 1}{p} - \frac{\rho - 1}{1-p} = 0 ,$$

and

$$n/\lambda - \sum_{i=1}^{n} \{N-p(i-1)\} t_i + \frac{\mu-1}{\lambda} - \gamma = 0.$$

Finally, the $100 \, (1-\alpha) \, \$$ Bayesian confidence region is given by

$$f(N,p,\lambda) = C$$

where

$$f(N,p,\lambda) = n\log\lambda - \sum_{i=1}^{n} \{N-p(i-1)\}t_{i}$$

$$+ \sum_{i=1}^{n} y_{i}\log(1-p) + \sum_{i=1}^{n} (1-y_{i})\log\{N-(i-1)\}$$

$$+ (\alpha-1)\log N - \beta N$$

$$+ (\mu-1)\log \lambda - \gamma \lambda$$

$$+ (\pi-1)\log p + (\rho-1)\log(1-p)$$

and

$$C = f(\hat{N}, \hat{p}, \hat{\lambda}) - \frac{1}{2} \chi_{3;\alpha}^{2}$$
.

5. BAYESIAN SOFTWARE CORRECTION LIMIT POLICIES

The objective of the investigation was to provide an optimum correction limit policy for a large-scale software system subject to random error occurrences and error removals in an operational phase. When an error occurs a corrective action is undertaken to remove it. Such an action can be scheduled at two levels, which we call Phase I and Phase II. By Phase I we mean that the corrective action will be undertaken by the programmer while Phase II action is undertaken by a system analyst or system designer. First, Phase I corrective action is scheduled for a specified time T . If the error is not corrected in this time, it is referred to Phase II. This sequence of corrective actions in an operational phase is shown in Figure 5.1. Our objective is to determine the optimum value T* of T which minimizes the long run average cost. Two models are developed for this purpose. In the first model we assume that the cost of observations of error occurrence and correction time, prior to the implementation of the optimum policy, is negligible. The second model incorporates the cost of observations.

Details of the model and related results are given in reference [8].

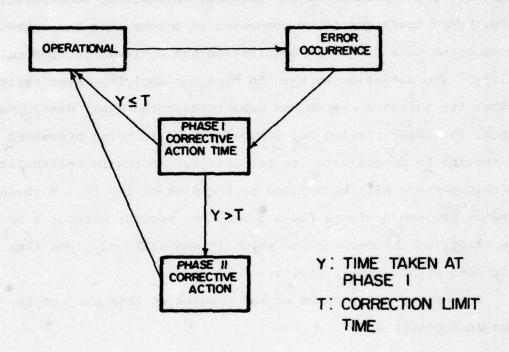


Figure 5.1 Sequence of Corrective Actions in Operational Phase

6. SOFTWARE RELIABILITY DEMONSTRATION TEST PLANS

The purpose of this task was to describe the theory, methodology, and procedures for software reliability demonstration tests. Such tests are to be conducted to ensure that the software system being considered for acquisition has achieved desired reliability. The situation we have in mind for applying these tests is that the software system has gone through the usual development phases, including testing and debugging, and is being presented for testing to demonstrate its reliability. Software reliability for this purpose will be defined as the probability that a given software program operates for a given time period, without a software error, on the machine for which it was designed given that it is used within design limits.

A complete description of the results on this project is given in Appendix B.

7. COMPUTER PROGRAMS

The computer programs required for computations of the quantities described in Sections 2 through 6 are given in Appendix C. The programs are self-explanatory, give the list of input/output parameters and include listings of the needed subroutines.

SELECTED REFERENCES

- [1] Akiyama, F. (1971), "An Example of Software Debugging," 1971 IFIP Congress, pp. TA-3-37 to TA-3-42.
- [2] Baker, W. F. (1977), "Software Data Collection and Analysis:

 A Real-Time System Project History," IBM Corporation,

 Final Technical Report, RADC-TR-77-192, June 1977.(A041644)
- [3] Belady, L. A. and Lehman, M. M. (1976), A Model of Large

 Program Development, IBM Systems Journal, Vol. 15, no. 3

 pp. 225-252.
- [4] Boehm, B. W., Brown, J. R., and Lipow, M. (1976), "Quantitative Evaluation of Software Quality," Proc. 2nd International Conference on Software Engineering, pp. 54-59.
- [5] Endres, A. (1975), "An Analysis of Errors and Their Causes in System Programs," Proceedings: 1975 International Conference on Reliable Software, pp. 327-336.
- [6] Fries, M. J. (1977), "Software Error Data Acquisition," Boeing Aerospace Company, <u>Final Technical Report</u>, RADC-TR-77-130, April 1977.(A039916)
- [7] Goel, A. L. and Okumoto, K. (1978), "An Imperfect Debugging Model for Reliability and Other Quantitative Measures of Software Systems," <u>Technical Report</u> No. 78-1, Department of IE & OR, Syracuse University.
- [8] Goel, A. L. and Okumoto, K. (1978), "Bayesian Software Correction Limit Policies," <u>Technical Report No. 78-4</u>, Department of IE & OR, Syracuse University.
- [9] Goel, A. L. and Okumoto, K. (1978), A Time Dependent Error Detection Rate Model for a Large-scale Software System," to appear in the <u>Third USA-Japan Computer Conference</u> Proceedings.
- [10] Halstead, M. H. (1977), Elements of Software Science, Elsevier.
- [11] Jelinski, J. and Moranda, P. B. (1972), "Software Reliability Research," in <u>Statistical Computer Performance Evaluation</u>, edited by W. Freiberger, Academic Press.
- [12] Jelinski, J. and Moranda, P. B. (1973), "Applications of a Probability-Based Model to a Code Reading Experiment,"

 Record: 1973 IEEE Symposium on Computer Software Reliability, pp. 78-81.

- [13] Jones, T. C. (1978), "Measuring Programming Quality and Productivity," IBM Systems Journal, Vol. 17, No. 1, pp. 39-63.
- [14] Littlewood, B. (1975), "A Reliability Model for Markov Structured Software," Applied Statist., Vol. 24, pp. 172-177.
- [15] Littlewood, B. and Verrall, J. L. (1973), "A Bayesian Reliability Growth Model for Computer Software," <u>Applied</u> <u>Statist.</u>, Vol. 22, pp. 332-346.
- [16] Miyamoto, I. (1975), "Software Reliability in On-Line Real Time Environment," <u>Proceedings: 1975 International Con-</u> ference on Reliable Software, pp. 194-203.
- [17] Moranda, P. (1975), "A Comparison of Software Error-rate Models," 1975 Texas Conference on Computing.
- [18] Motley, R. W. et al (1977), "Statistical Prediction of Programming Errors," IBM Corporation, Final Technical

 Report, RADC-TR-77-175, May 1977.(A041106)
- [19] Musa, J. D. (1975), "A Theory of Software Reliability and its Application," <u>IEEE Trans. on Software Engineering</u>, Vol. SE-1, No. 3, pp. 312-327.
- [20] Okumoto, K. and Goel, A. L. (1978), "Classical and Bayesian Inference for the Software Imperfect Debugging Model,"

 Technical Report No. 78-2, Department of IE & OR,
 Syracuse University.
- [21] Okumoto, K. and Goel, A. L. (1978), "Availability Analysis of Software Systems under Imperfect Maintenance," <u>Technical Report No. 78-3, Department of IE & OR,</u> Syracuse University.
- [22] Ross, S. M. (1970), Applied Probability Models With Optimization Applications, Holden-Day.
- [23] Roussas, G. G. (1973), A First Course in Mathematical Statistics, Addison-Wesley.
- [24] Rye, P. et al (1977), "Software Systems Development: A CSDL Project History," The Charles Stark Draper Laboratory, Inc., Final Technical Report, RADC-TR-77-213, June 1977. (A042186)
- [25] Schick, G. J. and Wolverton, R. W. (1972), "Assessment of Software Reliability," TRW Systems Group, <u>TRW Software</u> Series, TRW-SS-72-04, September 1972.

- [26] Schneidewind, N. J. (1972), "An Approach to Software Reliability Prediction and Quality Control," <u>AFIPS Conference</u> <u>Proceedings, Vol. 41 Part II</u>, Fall Joint Computer Conference, pp. 837-838.
- [27] Schneidewind, N. J. (1975), "Analysis of Error Processes in Computer Software," <u>Proceedings: 1975 International Con-</u> ference on Reliable Software, pp. 337-346.
- [28] Shooman, M. L. (1972), "Probabilistic Models for Software

 Reliability Prediction," Statistical Computer Perfor
 mance Evaluation, pp. 485-502, Academic Press, New York.
- [29] Shooman, M. L. and Bolsky, M. I. (1975), "Types, Distribution, and Test Correction Times for Programming Errors,"

 Proceedings: 1975 International Conference on Reliable Software, pp. 347-357.
- [30] Sukert, A. N. (1977), "An Investigation of Software Reliability Models," Proc. 1977 R & M.
- [31] Sukert, A. N. (1977), "A Multi-Project Comparison of Software Reliability Models," <u>Computers in Aerospace Conference</u>, pp. 413-421.
- [32] Thayer, T. A. et al (1976), "Software Reliability Study,"

 TRW Defense & Space Systems Group, Final Technical

 Report, RADC-TR-76-238, Aug. 1976.(A030798)
- [33] Trivedi, A. K. and Shooman, M. L. (1975), "A Many-State Markov Model for the Estimation and Prediction of Computer Software Performance Parameters," Proceedings: 1975 Inter-national Conference on Reliable Software, pp. 208-220.
- [34] Wagoner, W. L. (1973), "The Final Report on Software Reliability Measurement Study," <u>Aerospace Report No.</u> TOR-0074(4112)-1.
- [35] Willman, H. E. Jr. et al (1977), "Software Systems Reliability: A Raytheon Project History," Raytheon Company, <u>Final</u> <u>Technical Report</u>, RADC-TR-77-188, June 1977.(A040992)

APPENDIX A

AN ANALYSIS OF RECURRENT SOFTWARE ERRORS IN A REAL-TIME CONTROL SYSTEM

In this Appendix we present an analysis of software error data from a large-scale software project using the imperfect debugging model discussed in Section 2. The model parameters are estimated from the data and the values predicted from the model are compared with the observed values. Joint confidence regions for the parameters are also constructed which permit a study of the sensitivity of predictions.

A.1 Analysis of Error Data from a Real Time Control System

A real-time control system for a land-based radar system was developed by Raytheon Co. in a modular fashion. Nearly all of the modules were written in JOVIAL/J3. The error data base was extracted from 2165 Software Problem Reports (SPRs) written against 109 operational software modules over the development phases and is described in [35]. Table A.1 shows the distribution of the SPRs by month opened during a 22 month period of integration, acceptance and operational testing phases.

The available data give the number of total errors (u_i) and the number of imperfect debugging or recurrent errors (v_i) detected by time t_i . The parameters under consideration are the initial number of software errors (N), the error occurrence rate for each error (λ) , and the probability (p) of perfect debugging. We first estimate the parameters N, p and λ from these data (t,u,v) and then compare the results obtained from IDM with the observed values.

TABLE A.1
ERROR DATA BY MONTH

Month (t _i)	Total Number of Errors		Imperfect Debugging Errors	
	u _i - u _{i-1}	ui	v _i - v _{i-1}	v _i
1	122	122	6	6
2	98	220	1	7
3	82	302	1	8
4	75	377	3	11
5	113	490	2	13
6	85	575	3	16
7	105	680	2	18
8	47	727	1 1	19
9	61	788	1	20
10	25	813	1	21
11	28	841	0	21
12	42	883	1	22
13	18	901	0	22
14	17	918	0	22
15	28	946	0	22
16	14	960	0	22
17	5	965	0	22
18	3	968	0	22
19	3	971	0	22
20	13	984	0	22
21	5	989	0	22
22	10	999	0	22

Using the data in Table A.1 and the results from Section 2, the results are obtained as shown in Table A.2.

The joint confidence regions for N and λ for p=0.974 are plotted in Figure A.1. The plots of the actual and fitted SPR's by month are shown in Figure A.2, and the plots of actual and predicted number of remaining errors are given in Figure A.3.

TABLE A.2
A SUMMARY OF ERROR DATA ANALYSES

Quantities of Interest		Calculated Values	
û		1079 (Errors)	
ĝ		0.974	
î		0.1235 (per month)	
â (≡ <u>Ñ</u>)		1108 (Errors)	
Ĝ (≡ĝλ̂)		0.1203 (per month)	
	$\{\mathbf{N},\overline{\mathbf{N}}\}$	{914, 1244}	
sp	{ <u>p</u> , p }	(0.964, 0.984)	
spunoq	$\{\underline{\lambda},\overline{\lambda}\}$	{0.096, 0.151}	
1 %06	(<u>a</u> , a)	[933, 1283]	
6	(<u>b</u> , <u>b</u>)	{0.0926, 0.1480}	
[₽] Ñ, î		-0.723	
^p â,b		-0.743	

RAYTHEON 4000 = 1079 = 0.974 = 0.1235 3000 95% PARAMETER N 2000 (â,â) 75% 1000 0 0.1 0.2 0.3 0. 0.4 PARAMETER A

Figure A.1 Joint confidence regions for N and λ for $p\!\!=\!\!\hat{p}$

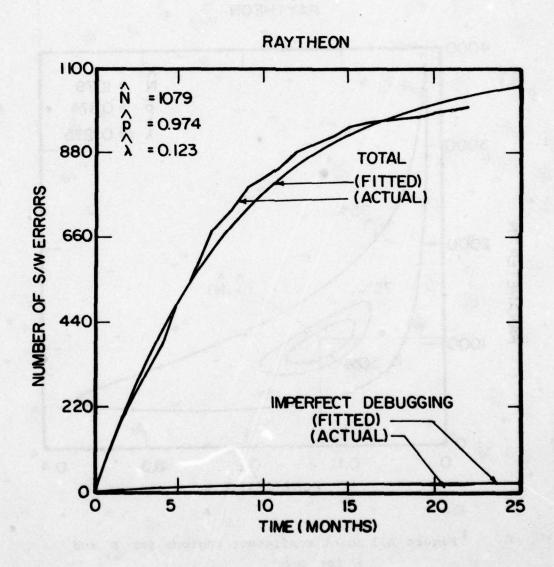


Figure A.2 Actual and fitted SPRs by month

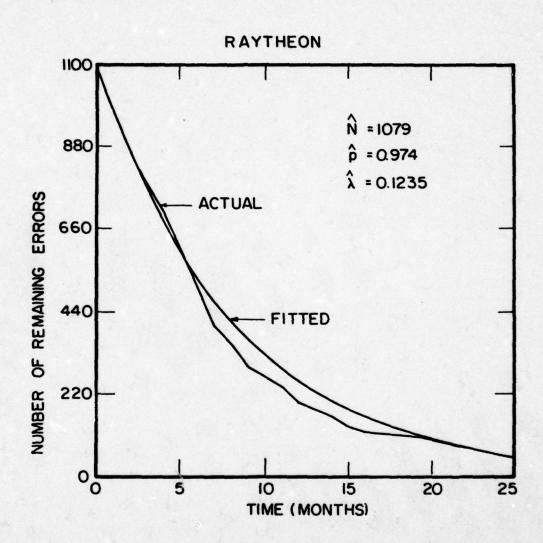


Figure A.3 Plots of the actual and predicted number of remaining errors.

APPENDIX B

SOFTWARE RELIABILITY DEMONSTRATION TEST PLANS

B.1. INTRODUCTION AND PROBLEM DEFINITION

The purpose of this report is to describe the theory, methodology, and procedures for software reliability demonstration tests. Such tests are to be conducted to ensure that the software system being considered for acquisition has achieved desired reliability. The situation we have in mind for applying these tests is that the software system has gone through the usual development phases, including testing and debugging, and is being presented for testing to demonstrate its reliability. Software reliability for this purpose will be defined as the probability that a given software program operates for a given time period, without a software error, on the machine for which it was designed, given that it is used within design limits.

In order to develop an appropriate test plan, we must first choose a model which adequately describes the error occurrence phenomenon. Most software error models are based on the assumption that successive errors follow a decreasing failure rate because of the reduction in the number of remaining errors. However, in this appendix we assume that the times between errors during the demonstration phase follow an exponential distribution, i.e. have

a constant failure rate. As a justification for this assumption, we contend that the demonstration time and the number of errors encountered during demonstration will be relatively small and hence a constant failure rate model will be an adequate representation of the error phenomenon. At worst, this model will yield a somewhat conservative test from the viewpoint of DOD.

Let the distribution of error occurrence times be given by

$$f(t|\lambda) = \lambda \exp(-t\lambda), \quad t \ge 0, \lambda > 0 \quad (B.1-1)$$

Then the following measures of software reliability can be used interchangeably:

. Reliability,
$$R(\tau) = P(t > \tau) = e^{-\tau \lambda}$$
 (B.1-2)

- . Meantime Between Software Failure (MTBSF) = $1/\lambda$
- . Software Failure Rate (SFR) = λ

In the following we shall take λ as the measure of software reliability.

In the classical set up, the demonstration tests are generally designed to distinguish between two values of λ , namely the maximum acceptable SFR, λ_1 and the specified SFR, λ_0 . The decision to accept or reject the software system is based upon test results which are subject to random fluctuation. A loss is incurred when either an accept or a reject decision is wrongly taken.

The loss corresponding to wrong decisions is quantified by two risks called the producer's risk, $\alpha = P(R|\lambda_0)$, and the consumer's risk $\beta = P(A|\lambda_1)$, where A and B denote acceptance and rejection of the

software system, respectively.

In this context various types of demonstration plans can be designed to limit α and β to desired values. For example, a truncated single sample plan for the system is employed as follows. The plan consists of using the software in an environment which is representative of the operational environment for a time period T. The number of errors r encountered during this time period is recorded. (In this study we assume that all errors are of the same severity. When errors are classified according to the degree of severity, the problem becomes quite complicated and is beyond the scope of this investigation). If r is less than or equal to a pre-specified number, r*, the software is accepted. Otherwise, the system is rejected. The design of such a plan consists of obtaining the quantities T and r* such that the desired risks α and β are satisfied.

One disadvantage in the above approach is that λ is assumed to be an unknown, fixed constant. In practice there may be sufficient reason to believe that knowledge about λ is available in some quantifiable form and one would like to incorporate such information in the development of the demonstration test plan. Another important situation arises when the risks associated with the demonstration test are not adequately represented by the classical risks (α,β) and interest lies in associating risks with the posterior distribution of λ . In such cases one resorts to a Bayesian approach for test design.

The problem under consderation then, is the design of single sample software reliability demonstration plans when error occurrence times are assumed to follow the exponential distribution.

Before describing the development of the test plans, we first discuss the various risks that arise in the above situations in Section B.2. Fixed time, classical tests are then considered in Section B.3. In Sections B.4 and B.5 Bayesian tests are developed for two situations: (i) λ has a noninformative prior distribution, and (ii) information about λ can be quantified from the testing and debugging data.

Section B.6 presents the step-by-step procedure to be employed to demonstrate the properties and performance of the demonstration test plans.

B. 2. DEFINITIONS AND INTERPRETATION OF RISKS

The following quantities are of interest

- 1. The probability $P(R|\lambda = \lambda_0) = \alpha$ that a software system with specified SFR is rejected.
- 2. The probability $P(A|\lambda = \lambda_1) = \beta$ that a software system with maximum acceptable SFR is accepted.
- 3. The probability $P(A|\lambda \ge \lambda_1) = \overline{\beta}$ that a software system which is of unacceptable reliability is accepted.
- 4. The probability $P(R|\lambda \le \lambda_0) = \bar{\alpha}$ that a software system of acceptable reliability is rejected.
- 5. The probability $P(\lambda \ge \lambda_1 | A) = \beta *$ that the SFR of a system which has been accepted is more than the maximum acceptable SFR.
- 6. The probability $P(\lambda \ge \lambda_0 | A) = \beta^{**}$ that the SFR of a system which has been accepted is more than the specified SFR.
- 7. The probability $P(\lambda \le \lambda_0 | R) = \alpha^*$ that the SFR of a system which has been rejected is less than the specified SFR.
- 8. The probability P(R) that the software system is rejected.
- 9. The probability P(A) that the software system is accepted.
- 10. The probability $P(\lambda \ge \lambda')$ that a-priori, the software system has SFR which is more than λ' .

These quantities will now be discussed in some detail.

B. 2.1, Classical Risks (α,β)

The classical producer's risk α and consumer's risk β are defined as follows:

- $\alpha = P(R|\lambda = \lambda_0)$, the probability of rejecting a (B.2-1) software system whose SFR is equal to the specified value, λ_0 .
- β = P(A| λ = λ_1), the probability of accepting a software (B.2-2) system whose SFR is equal to the maximum acceptable value.

The (α,β) risks represent two points on the classical operating characteristic (OC) curve which is a plot of $P(A|\lambda)$ versus λ . These risks do not provide an explicit control of the probability of acceptance for values of λ other than λ_1 and λ_0 . However, $P(A|\lambda)$ decreases monotonically with λ . Hence, if $\lambda < \lambda_0$, the probability of rejection is less than α . If $\lambda > \lambda_1$, the probability of acceptance is less than β . The shape of the OC curve governs the degree of protection provided in the indifference zone between λ_1 and λ_0 .

B. 2.2. Average Risks (α, β)

The average risks are defined as follows

- $\bar{\alpha}$ = P(R| $\lambda \le \lambda_0$), the probability of rejecting a software (B.2-3) system with a SFR less than or equal to the specified value, λ_0 .
- $\overline{\beta}$ = P(A| $\lambda \ge \lambda_1$), the probability of accepting a software (B.2-4) system with a SFR greater than or equal to λ_1 .

Mathematically, the risks may be expressed as:

$$P(R|\lambda \leq \lambda_0) = \frac{P(R,\lambda \leq \lambda_0)}{P(\lambda \leq \lambda_0)} = \frac{\int_0^{\lambda_0} P(R|\lambda)p(\lambda)d\lambda}{\int_0^{\lambda_0} p(\lambda)d\lambda}$$
(B.2-5)

$$\bar{\alpha} = \int_{0}^{\lambda_{0}} P(R|\lambda) \cdot p(\lambda|\lambda \le \lambda_{0}) d\lambda$$
 (B.2-6)

and, similarly

$$\bar{\beta} = \int_{\lambda_1}^{\infty} P(A|\lambda) \cdot P(\lambda|\lambda \ge \lambda_1) d\lambda$$
 (B.2-7)

The average risks provide the following protection. If the producer produces a large number of software systems, then, in the long run, less than $\bar{\alpha}$ percent of the desired ones will be rejected. If the consumer buys a large number of software systems, then, in the long run, less than 100 $\bar{\beta}$ percent of the bad systems will be accepted. No explicit control on the probability of acceptance is provided at any specific value of λ when we use the average risk criteria.

2.3 Posterior Risks (α*, β*)

The (α^*, β^*) risks are defined as follows

$$\alpha^* = P(\lambda \leq \lambda_0 | R)$$

This risk is the long run probability of a rejected software system being good.

$$\beta^* = P(\lambda \ge \lambda_1 | A)$$

This risk is the long run probability of an accepted system being bad. Mathematically,

$$\alpha^* = P(\lambda \le \lambda_0 | R) = \int_0^{\lambda_0} P(\lambda | R) d\lambda = \frac{\int_0^{\infty} P(R | \lambda) p(\lambda) d\lambda}{\int_0^{\infty} P(R | \lambda) p(\lambda) d\lambda}$$

$$\beta * = P(\lambda \ge \lambda_1 | A) = \begin{cases} \int_{0}^{\infty} P(A | \lambda) p(\lambda) d\lambda \\ \frac{\lambda_1}{\int_{0}^{\infty} P(A | \lambda) p(\lambda) d\lambda} \end{cases}$$

These risks can also be interpreted in a "degree of belief" sense. Thus, α^* would represent a persons's degree of belief that if a software system has been rejected, it has a SFR which is better than the specified value. Similarly, β^* would be the degree of belief that the SFR of an accepted system is worse than the maximum acceptable value.

B. 2.4 Probability of Rejection P(R)

This is a single number given by

$$P(R) = \int_{0}^{\infty} P(R|\lambda)p(\lambda)d\lambda \qquad (B.2-12)$$

or

$$P(R) = 1 - P(A) = 1 - \int_{0}^{\infty} P(A|\lambda)p(\lambda)d\lambda \qquad (B.2-13)$$

Note that the integration is over the entire range of λ and specification of λ_0 , which is usually specified in conjunction with a producer's risk, is unnecessary.

In the frequency sense we have

For the producer this criterion implies that, in the long run, less than (100) P(R) percent of the systems will be rejected.

B.2.5. Alternate Posterior Consumer's Risk g**

We define a new risk associated with the posterior distribution of λ as follows

$$\beta^{\star\star} = P(\lambda \ge \lambda_0 \big| A) = \int\limits_{\lambda_0}^{\infty} f(\lambda \big| A) d\lambda, \qquad (B. 2-14)$$
 Where $f(\lambda \big| A)$ is the pdf of λ conditional on acceptance. This can be

written as

$$\beta^{**} = \frac{\int_{0}^{\infty} P(A|\lambda) p(\lambda)d\lambda}{\int_{0}^{\infty} P(A|\lambda) p(\lambda)d\lambda}$$
(B. 2-15)

This risk gives the long run probability of the accepted system having a λ above the specified SFR λ_0 .

B. 3. DESIGN OF TEST PLAN FOR CLASSICAL RISKS

Now we consider the design of a single sample plan where the parameters are T and r*. Since the time to software failure is exponential, the observed number of software errors r in fixed time T has a Poisson distribution, i.e.

$$f(r|\lambda) = \frac{e^{-T\lambda} (T\lambda)^r}{r!}$$
 (B.3-1)

The two risks can be written as:

$$\sum_{r=0}^{r^*} \frac{e^{-T\lambda_0} (T\lambda_0)^r}{r!} = 1 - \alpha$$
 (B.3-2)

and

$$\sum_{r=0}^{r^*} \frac{e^{-T\lambda_1} (T\lambda_1)^r}{r!} = \beta$$
 (B.3-3)

Given λ_1 , λ_0 , α and β , the above equations can be simultaneously solved to obtain software test time T and allowable number of errors r*.

Individual specification of λ_1 and λ_0 is not necessary. Let $K = \lambda_1/\lambda_0$, and let $T^* = T \lambda_0$. Then, to satisfy the stated risks we can write

$$\sum_{r=0}^{r*} \frac{e^{-T*} (T*)^r}{r!} \ge 1 - \alpha$$
 (B.3-4)

and

$$\sum_{r=0}^{r*} \frac{e^{-KT*} (KT*)^r}{r!} \leq \beta$$
 (B.3-5)

Given (K, α, β) the above two equations can be solved to obtain (T^*, r^*) . These equations can be solved numerically or by using tables of cumulative Poisson probabilities. Clearly, test plans with identical λ_0/λ_1 have the same T^* and r^* values. Given the specified value λ_0 , the actual test time is obtained as $T = T^*/\lambda_0$.

4. BAYESIAN TEST PLANS FOR NONINFORMATIVE PRIOR DISTRIBUTION

In this section we consider the case when it is possible to express the unknown parameter λ in terms of a metric $\phi(\lambda)$ so that the corresponding likelihood is data translated. This means that the likelihood function for $\phi(\lambda)$ is completely determined a-priori except for its location which depends on the software failure data yet to be observed. This state of indifference can be expressed by taking $\phi(\lambda)$ to be locally uniform, and the resulting prior distribution is called noninformative for $\phi(\lambda)$ with resepect to data. A more detailed discussion on noninformative prior distribution tan be found in Box and Tiao [1]* . In our case, a noninformative prior for λ is

$$p(\lambda) \propto \lambda^{-1/2}$$
 (B.4-1)

Now, if T is the test time for a software system, the number of errors in T will be given by a Poisson distribution with parameter $T\lambda$ as mentioned earlier. Letting Λ = $T\lambda$, the prior for Λ can be written as

$$P(\Lambda) \propto \Lambda^{-1/2} \tag{B.4-2}$$

^{*}References at the end of this Appendix

The joint probability distribution of Λ and r is:

$$p(r, \Lambda) = p(r|\Lambda) \cdot p(\Lambda)$$

$$\propto \frac{(\Lambda)^{r} \cdot e^{-\Lambda}}{r!} \cdot \Lambda^{-1/2}$$

or

$$p(r, \Lambda) = b' \frac{\Lambda^{r-1/2} \cdot e^{-\Lambda}}{r!}$$
 (B.4-3)

where b' is the normalizing constant.

To get the marginal distribution of r, we let Λ be defined over the range (0,d) where d is sufficiently large and d < ∞ . Then we have

$$p(r) = \int_{0}^{\infty} P(r, \Lambda) \cdot d\Lambda = \int_{0}^{\infty} b^{r} \cdot \frac{\Lambda^{r-1/2} \cdot e^{-\Lambda}}{r!} d\Lambda \qquad (B4-4)$$

Now we choose d such that

$$\int_{0}^{\infty} \frac{b' \Lambda^{r-1/2} \cdot e^{-\Lambda}}{r!} d\Lambda < \varepsilon$$
 (B4-5)

for some given sufficiently small $\varepsilon > 0$: Then

$$p(r) = \frac{b!}{r!} \Gamma(r + \frac{1}{2})$$
 (R.4-6)

and

$$P(A) = \sum_{r=0}^{r*} p(r) = \sum_{r=0}^{r*} \frac{b'\Gamma(r + \frac{1}{2})}{r!}$$
 (B(4-7)

In order to get expressions for various risks, we proceed as follows.

Let $\Lambda_0 = T\lambda_0$. Then from (B.4-3) we get,

$$P(\lambda \le \lambda_0, \text{ accept}) = P(\Lambda \le \Lambda_0, \text{ accept})$$

$$= \sum_{r=0}^{r*} \int_{0}^{\Lambda_0} b' \cdot \frac{\Lambda^{r-1/2} \cdot e^{-\Lambda}}{r!} d\Lambda \qquad (B.4-8)$$

Substituting the above expressions into the appropriate formulae in Section 2, we get the equations for desired risks. For example, if we are interested in the risk combination $(\bar{\alpha}, \beta^*)$, we get from (B.2-3) and (B.2-9), respectively:

$$\bar{\alpha} = 1 - P(A | \lambda \le \lambda_0)$$

$$= 1 - \frac{P(A, \Lambda \le \lambda_0)}{P(\lambda \le \lambda_0)}$$

$$= 1 - \frac{\sum_{r=0}^{r+1} \int_0^{\Lambda_0} b' \cdot \frac{\Lambda^{r-1/2} \cdot e^{-\Lambda}}{r!} d\Lambda$$

$$= 1 - \frac{\Lambda_0}{\Lambda_0}$$

$$\int_0^{\Lambda_0} b' \Lambda^{-1/2} d\Lambda$$

or

$$\bar{\alpha} = 1 - \frac{\sum_{r=0}^{r^*} \int_{0}^{\Lambda_0} \frac{\Lambda^{r-1/2} \cdot e^{-\Lambda}}{r!} d\Lambda}{\int_{0}^{\Lambda_0} \Lambda^{-1/2} d\Lambda}$$
(B.4-9)

and

$$\beta^* = P(\lambda \ge \lambda_1 | A)$$

$$= 1 - P(\lambda \le \lambda_1 | A)$$

$$= 1 - \frac{\sum_{r=0}^{r^*} \int_0^{\Lambda_1} b' \Lambda^{r-1/2} \cdot e^{-\Lambda} d\Lambda}{\sum_{r=0}^{r^*} \frac{b' \Gamma(r + \frac{1}{2})}{r!}}$$

or

$$\beta^* = 1 - \frac{\sum_{r=0}^{r^*} \int_{0}^{\Lambda_1} \Lambda^{r-1/2} \cdot e^{-\Lambda} d\Lambda}{\sum_{r=0}^{r^*} \frac{\Gamma(r + \frac{1}{2})}{r!}}$$
(B.4-10)

If the consumer's risk of interest is $\beta **$, then from (B.2-14) we have

$$\beta^{**} = P(\lambda \ge \lambda_0 | A)$$

$$= 1 - P(\lambda \le \lambda_0 | A)$$

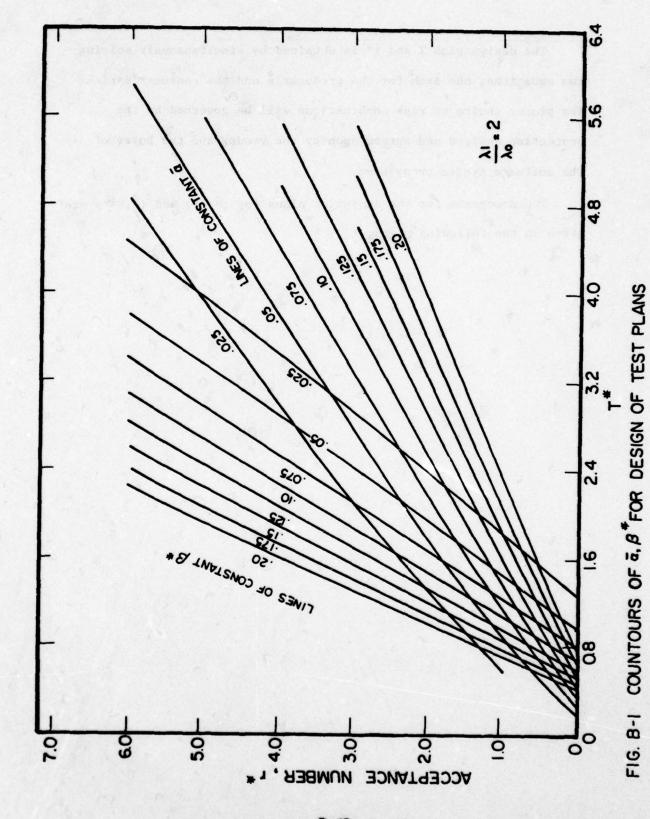
$$= \frac{\sum_{r=0}^{r^*} \int_{0}^{\Lambda_0} \Lambda^{r-1/2} \cdot e^{-\Lambda} d\Lambda}{\sum_{r=0}^{r^*} \frac{\Gamma(r + \frac{1}{2})}{r!}}$$
(B. 4-11)

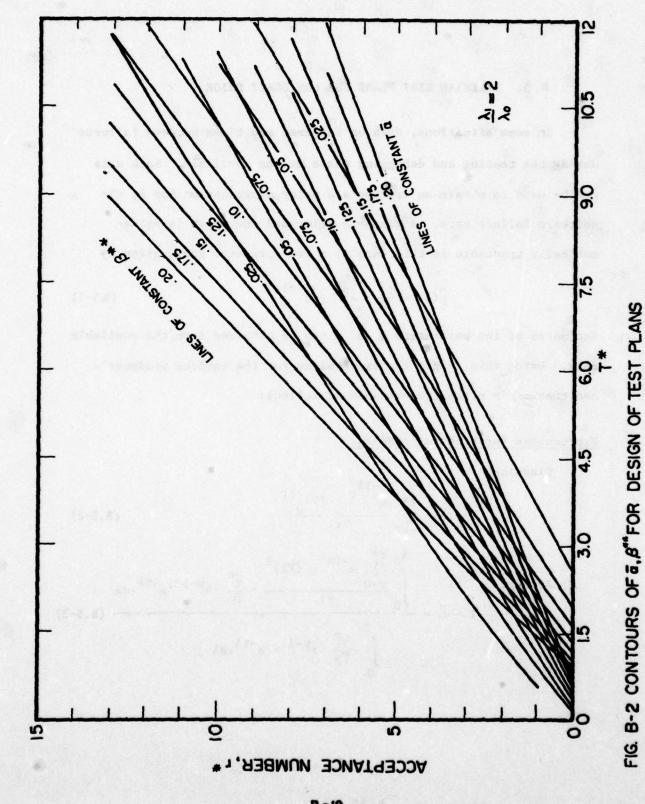
B- 16

The design plan T and r* is obtained by simultaneously solving two equations, one each for the producer's and the consumer's risk. The proper choice of risk combinations will be governed by the protection desired and agreed upon by the vendor and the buyer of the software system or systems.

Two nomograms for the design of plans for $(\bar{\alpha}, \beta *)$ and $(\bar{\alpha}, \beta **)$ are given on the following pages,







B. 5. BAYESIAN TEST PLANS FOR CONJUGATE PRIOR

In some situations, data on failures and times between failures during the testing and debugging phase may be available. Such data can be used to obtain an appropriate prior distribution for λ , the software failure rate. A flexible prior and one which is mathematically tractable in this case is a two parameter gamma given by

$$p(\lambda) = \frac{\tau^{p}}{\Gamma p} \cdot \lambda^{pT} \cdot e^{-\tau \lambda}$$
 (B.5-1)

Estimates of the parameters p and τ can be obtained from the available data. Using this prior, the expressions for the various producer's and consumer's risks are obtained as follows:

Expressions for Producer's Risks:

Classical:

Average:
$$1 - \alpha = \sum_{r=0}^{r^*} \frac{e^{-T\lambda_0} \cdot (T\lambda_0)^r}{r!} \qquad (B.5-2)$$

$$\frac{1}{1 - \alpha} = \frac{\int_0^0 \frac{\sum_{r=0}^{r^*} e^{-T\lambda} \cdot (T\lambda)^2}{\sum_{r=0}^{r^*} r!} \cdot \frac{\tau^p}{\overline{lp}} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}{\int_0^1 \frac{\tau^p}{\overline{lp}} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda} \qquad (B.5-3)$$

BOSCHERME WARRES

Posterior:

$$1 - \alpha \star = \frac{\int_{0}^{\lambda_{0}} \left\{1 - \sum_{r=0}^{r \star} \frac{e^{-T\lambda} \cdot (T\lambda)^{r}}{r!}\right\} \cdot \frac{\tau^{p}}{\Gamma p} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}{1 - \int_{0}^{\infty} \left\{\sum_{r=0}^{r \star} \frac{e^{-T\lambda} (T\lambda)^{r}}{r!}\right\} \cdot \frac{\tau^{p}}{\Gamma p} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}$$
(B. 5-4)

Probability of Rejection:

$$P(R) = 1 - \int_{0}^{\infty} \left\{ \sum_{r=0}^{r*} \frac{e^{-T\lambda} (T\lambda)^{r}}{r!} \right\} \cdot \frac{\tau^{p}}{\Gamma p} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda$$
 (B. 5-5)

Expressions for Consumer's Risks:

Classical:

$$\beta = \sum_{r=0}^{r*} \frac{e^{-TK\lambda_0} \cdot (KT\lambda_0)^r}{r!}$$
(B. 5-6)

Average:

$$\bar{\beta} = \frac{\int_{\lambda_{1}}^{\left\{\sum_{r=0}^{r^{*}} \frac{e^{-T\lambda} \cdot (T\lambda)^{r}}{r!}\right\} \cdot \frac{\tau^{p}}{\Gamma_{p}} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}{\int_{\lambda_{1}}^{\infty} \frac{\tau^{p}}{\Gamma_{p}} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d^{\lambda}}$$
(B. 5-7)

Posterior:

$$\beta^{\star} = \frac{\int_{\lambda_{1}}^{\infty} \left\{ \sum_{r=0}^{r^{\star}} \frac{e^{-T\lambda} \cdot (T\lambda)^{r}}{r!} \right\} \frac{\tau^{p}}{\Gamma_{p}} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}{\int_{0}^{\infty} \left\{ \sum_{r=0}^{r^{\star}} \frac{e^{-T\lambda} \cdot (T\lambda)^{r}}{r!} \right\} \frac{\tau^{p}}{\Gamma_{p}} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}$$
(B.5-8)

Alternate Posterior:

$$\beta^{**} = \frac{\int_{\lambda_0}^{\infty} \left\{ \sum_{r=0}^{r^*} \frac{e^{-T\lambda} (T\lambda)^r}{r!} \right\} \frac{\tau^p}{\Gamma p} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}{\int_{0}^{\infty} \left\{ \sum_{r=0}^{r^*} \frac{e^{-T\lambda} \cdot (T\lambda)^r}{r!} \right\} \frac{\tau^p}{\Gamma p} \cdot \lambda^{p-1} \cdot e^{-\tau\lambda} \cdot d\lambda}$$
(B.5-9)

B. 6. PROCEDURE FOR DEMONSTRATING THE PERFORMANCE OF THE TEST PLANS

This section provides a step-by-step procedure that will be used to demonstrate the properties and performance of the demonstration test plans for a software system as developed in the previous sections.

These procedures will be followed in conducting the demonstration on site.

Prior to conducting the demonstration test in practice, the DOD must make sure that the software system has been developed according to specifications and all the stated development phases have been carried out to meet the objectives.

The steps to be followed for a demonstration test are as follows.

- Decide the risk criteria (producer's and consumer's risks) for demonstration.
- 2. Choose the values of the two risks.
- Design a test plan (T,r*) which meets the risks in Step 2 as closely as possible. (Designs will be obtained using computer programs developed for this purpose.)
- 4. Simulate software errors using the software error simulator.
- Apply the demonstration test of Step (3) to simulated errors and make accept/reject decisions.
- 6. Compare the recorded results with theoretical results.

B.7. REFERENCES

- [1] Box, G.E.P. and Tiao, G.C. (1973)

 Bayesian Inference in Statistical Analysis. Addison-Wesley
- [2] Goel, A.L. and Joglekar, A.M. (1976)

"Reliability Acceptance Sampling Plans Based Upon Prior Distribution: Risk Criteria and Their Interpretation," RADC-TR-76-294, Volume II. (A033516)

APPENDIX C

COMPUTER PROGRAMS

C.1 Programs for the Imperfect Debugging Model (Section 2)

Computer programs to compute the following quantities, for given N, p and λ are given in Tables C.1 to C.8.

- Mean and variance of the first passage time from N to n_0 (SUBROUTINE FRSTRT)
- Pdf and cdf of the first passage time from N to n_0 (SUBROUTINE FPT2)
- Probability of the software system having n₀ errors remaining at time t (SUBROUTINE STATE)
- Expected number of errors remaining at time t and expected number of total errors and imperfect debugging errors detected by time t (SUBROUTINE MEAN)
- MTTF and reliability for given k at time x (SUBROUTINE TBF)

The programs are self-explanatory and include the required subroutines.

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY PARKISHED TO DDC

TABLE C.1 SUBROUTINE FRSTPT

```
subroutine frstpt(n,p,r,n0,et,vt,a,b)-----
                         - compute mean and variance of first passage time
   function
                               from n to nO , and estimate shape and scale
C
                               parameters of Gamma distribution
C
                         - call frstpt (n.p.r.nO.et.vt.a.b)
   usage
C
                         - (input.) initial no. of errors
C
   parameters
                      n
                         - (input.) prob. of perfect debussins
C
                      P
                         - (input.) detection rate / error - (input.) desired no. of errors
C
                     nO
C
                     et - (output.) mean time from n to no
C
                     vt - (output.) variance of time from n to n0
a - (output.) shape parameter of Gamma distribution
C
C
                         - (output.) scale parameter of Gamma distribution
C
      subroutine frstpt (n,p,r,n0,et,vt,a,b)
      x1=0.0
      x2=0.0
      n1=n0+1
      do 55 j=n1,n
      x1=x1+1.0/float(j)
      x2=x2+1.0/float(j)**2
   55 continue
      et=x1/p/r
      vt=x2/p/p/r/r
      b=et/vt
      a=et*b
      return
     , end
```

TABLE C.2 SUBROUTINE FPT2

```
subroutine fpt2 (n,p,r,n0,t,pdf,cdf)-----
C
C
   function
                        - compute p.d.f. and c.d.f. of first passage time from
                             n to no
C
C
                        - call fet2 (n.p.r.nO,t.edf,cdf)
   usage
C
   parameters
                      n - (input.) initial no. of errors
                      p - (input.) prob. of perfect debussins
C
                     r - (input.) detection rate / error
n0 - (input.) desired no. of errors
C
C
C
                      t - (input.) time
C
                    Pdf - (output.) p.d.f. of first passage time from n to no
                    cdf - (output.) c.d.f. of first passage time from n to n0
C
                        - frstpt, mdsamm
C
   read. subroutines
C
C
      subroutine fpt2 (n,p,r,n0,t,pdf,cdf)
      call frstpt (n,p,r,n0,et,vt,a,b)
      x=b*t
      x1=(a-1.0)*alog(x)
      x2=alsamma(a)
      x3=x1-x-x2
      if (x3 .1t. -88.0) so to 33
      Pdf=exp(x3)
                                               THIS PAGE IS BEST QUALITY PRACTICABLE
      so to 44
                                               FROM COPY PURMISHED TO DDC
   33 Pdf=0.0
   44 call mdsamm (x,a,cdf)
      return
   22 cdf=0.0
      Pdf=0.0
      return
      end
```

TABLE C.3 SUBROUTINE STATE

```
subroutine state (n,p,r,n0,t,pt)-----
   function - compute probability of being nO errors remaining
C
                            at time t
C
                      - call state (n.p.r.nO,t.pt)
C
   eksau
                     n - (input.) initial no. of errors
   parameters
C
                    p - (input.) prob. of perfect debussins
                    r - (input.) detection rate / error
n0 - (input.) desired no. of errors
t - (input.) time
C
C
C
                    pt - (output.) prob. of being nO errors at time t
   read. subroutines - fpt2
C
C
      subroutine state (n,p,r,n0,t,pt)
      if (n0.ea.n) so to 33
      call frt2(n.p.r.nO.t.pdf.cdf)
      pt1=cdf
      so to 44
   33 pt1=1.0
   44 n1=n0-1
      if (n1) 11,22,22
   11 pt2=0.0
     so to 99
   22 call fpt2 (n,p,r,n1,t,pdf,cdf)
      Pt2=cdf
   99 pt=pt1-pt2
      return
      end
```

TABLE C.4 SUBROUTINE MEAN

```
subroutine mean (n.p.r.t.er.ed0.ed1) -----
C
                         - compute expected no. of errors remaining at time to
   function
                              detected, and detected due to imperfect debussins
c
                              by time t
C
                         - call mean (n.p.r.t.er.ed0.ed1)
c
   ezezu
                     n - (input.) initial no. of errors
  parameters
C
                        - (input.) prob. of perfect debussins
C
                     -
                     r - (input.) detection rate/ error
t - (input.) time
C
C
                  er - (output.) expected no. of errors remaining at time t ed0 - (output.) expected no. of errors detected by time t ed1 - (output.) expected no. of imperfect debugging errors
C
C
C
C
                                           detected by time t
C
C
       subroutine mean (n,p,r,t,er,ed0,ed1)
       er=float(n)*exp(-p*r*t)
       ed0=(float(n)-er)/p
       ed1=ed0*(1.0-p)
       return
       end
```

TABLE C.5 SUBROUTINE TBF

```
subroutine tbf (n,p,r,k,x,rel,xmttf)-----
C
c function - compute reliability and mttf
c usase - call tbf (n,p,r,k,x,rel,xmttf)
c parameters n - (input.) initial no. of errors
                       p - (input.) prob. of perfect debussins
C
                      r - (input.) dec tion rate / error
C
                      k - (input.) k-th failure
x - (input.) time
C
C
                     rel - (output.) reliability at time x after (k-1)st failure
C
                   xmttf - (output.) mean time between (k-1)st and
C
                                         k-th failures
C
                                         ------
C
C
      subroutine tbf (n,p,r,k,x,rel,xmttf)
      xmttf=1.0/(float(n)-p*float(k-1))/r
    rel=exp(-x/xmttf)
      return
```

end

TABLE C.6 SUBROUTINE MDGAMM

```
subroutine mdsamm (x,p,prob)-----from imsl----
   function
                       - compute imcomplete samma distribution
C
   usade
                        - call mdsamm (x,p,prob)
C
   Parameters
                      x ~ (input.) value to which samma is to be integrated
C
                     P ~ (input.) gamma parameter
                  Prob - (output.) prob.=integral of samma(p) to x
   read. subroutines - samma
      subroutine mdsamm (x,p,prob)
      dimension v(6),v1(6)
      equivalence (v(3),v1(1))
      Prob=0.0
      if (x .se. 0.0) so to 5
      so to 9000
    5 if (p .st. 0.0) so to 10 so to 9000
   10 if (x .ea. 0.0) so to 9005
      fnls=alsamma(p)
      cnt=p*alog(x)
      went=x+fnls
      if ((cnt-yent) .st. -88.0) so to 15
      ax=0.0
      so to 20
   15 ax=exp(cnt-yent)
   20 bis=1.e35
      cut=1.e-8
      if ((x .le. 1.0) .or. (x .lt. p)) so to 40
      y=1.0-p
      z=x+y+1.0
      cnt=0.0
      v(1)=1.0
      v(2)=x
      v(3)=x+1.0
      v(4)=z*x
      Prob=v(3)/v(4)
   25 cnt=cnt+1.0
                                          THIS PAGE IS BEST QUALITY PRACTICABLE
      y=y+1.0
      z=z+2.0
                                            FROM COPY FURNISHED TO DDC
      went=y*ent
      v(5)=v1(1)*z-v(1)*yent
      v(6)=v1(2)*z-v(2)*yent
      if (v(6) .ea. 0.0) so to 50
      ratio=v(5)/v(6)
      reduc=abs(prob-ratio)
      if (reduc .st. cut) so to 30 if (reduc .le. ratio*cut) so to 35
```

```
30 prob=ratio
                            THIS PAGE IS BEST QUALITY PRACTICABLE
     so to 50
  35 prob=1.0-prob*ax
                               FROM COPY PURMISHED TO DDC
     so to 9005
  40 ratio=p
     cnt=1.0
     Prob=1.0
  45 ratio=ratio+1.0
     cnt=cnt*x/ratio
     prob=probtent
     if (cnt .st. cut) so to 45
     Prob=prob*ax/p
  so to 9005
50 do 55 i=1,4
     v(i)=v1(i)
  55 continue
     if (abs(v(5)) .lt. bis) so to 2 do 60 i=1,4
     v(i)=v(i)/bis
  60 continue
     so to 25
9000 continue
9005 return
     end
     function algamma(p)
     if (p .st. 31.0) so to 15
     call samma (P,SP)
     (ab)bols=smmsbls
     return
 15 z1=(p-0.5)*alog(p)-p+0.5*alog(2.0*3.1415)
    z2=1.0/12.0/P
     z3=1.0/360.0/P/P/P
     z4=1.0/1260.0/P/P/P/P/P
    z5=1.0/1680.0/p/p/p/p/p/p/p
     alsamma=z1+z2-z3+z4-z5
    return
    end
```

TABLE C.7 SUBROUTINE GAMMA

```
subroutine samma (xx,sx) -----
                                             -----from imsl--
   function
                      - compute a samma function of parameter xx
C
   usade
                      - call samma (xx, sx)
                  xx - (input.) parameter of samma function sx - (output.) value of samma function
C
   parameters
C
C
      subroutine samma(xx,sx)
      if(xx-57.) 6,6,4
    4 sx=1.0e30
      return
    xx=x
      err=1.e-6
      sx=1.
      if(x-2.) 50,50,15
   10 if (x-2.) 110,110,15
   15 x=x-1.
      XXXE=XX
   #o to 10
50 if (x-1.) 60,120,110
   60 if (x-err) 62,62,80
   62 y=float(int(x))-x
if (abs(y)-err) 120,120,64
   64 if (1.-w-err) 120,120,70
   70 if(x-1.) 80,80,110
   80 SX=SX/X
      x=x+1.
      so to 70
  110 y=x-1.
      ##=1.+##(-0.5771017+##(0.9858540+##(-0.8764218+##(0.8328212+##(-0.5684729
\c+u*(0.2548205+u*(-0.0514993)))))))
      EX=EXES
  120 return
      end
```

C.2 Programs for Simulation of Imperfect Debugging Data

Computer programs to simulate the data required to estimate the parameters N, p and λ , are given in Tables C.8 to C.11. These programs perform the following functions:

- Simulate data (t,y) for given N, p and λ (SUBROUTINE SMLT)
- Compute the mle's of N, p and λ given the data (t,χ) and also obtain the estimate of variance-covariance matrix
 (SUBROUTINE MLE)
- Compute the Bayesian estimates of N, p and λ for given data (t,y)(SUBROUTINE BAYES)

The programs are self-explanatory.

TABLE C.8 SUBROUTINE SMLT

```
subroutine smlt (n,p,r,nn,iseed,t,iy)----
C
   function
                    - simulate time between s/w failures for
                         imperfect debussins model
C
                    - call smlt (n.p.r.nn.iseed.t.iy)
   usade
C
                  n - (input.) initial no. of s/w errors
   parameters
                  p - (input.) prob. of perfect debussins
                  r - (input.) detection rate / error
              nn - (input.) no. of observations for s/w failure time
iseed - (input.) an integer value in the exclusive
                         rande (1,2147483647). iseed is replaced by
                          a new iseed to be used in subsequent calls.
                   t - output vector of length nn, containing time
                          between s/w failures
C
                 is - output vector of length nn, indicating the
c
                          type of error which is 1 if the i-th
                          failure is caused by an error due to
C
                          imperfect debussing, or O otherwith.
C
   read. subroutine - ssub
C
      subroutine smlt (n,p,r,nn,iseed,t,iy)
      dimension t(nn), iy(nn), rr(2)
      nr=n
      ie=0
      do 5 i=1,nn
                                        THIS PAGE IS BEST QUALITY PRACTICABLE
      if (nr .ea. 0) so to 99
      xm=1.0/r/float(nr)
                                       FROM COPY PARKISHED TO DDC
      call ssub (iseed,1,rr)
      t(1)=-xm*alog(rr(1))
      call ssub (iseed,1,rr)
      if (rr(1)-(1.0-p)) 22,22,33
   33 nr=nr-1
      so to 44
   22 ie=ie+1
   44 call ssub (iseed,1,rr)
      if (rr(1)-float(ie)/float(nr)) 55,55,66
   66 is(i)=0
      so to 77
   55 iy(i)=1
      ie=ie-1
   77 print 100, i, t(i), iy(i), nr, ie
  100 format(i5,f15.5,3i5)
    5 continue
   99 return
      end
```

TABLE C.9 SUBROUTINE MLE

```
subroutine mle (trigran, en, ep, er, ecov)-----
                  - estimate unknown parameters n, p, and lambda
C
                      and also estimate variance-covariance
                      matrix for idm
                  - call mle (t,iy,nn,en,ep,er,ecov)
C
                t - (input.) a vector of length nn, containing time
C
   parameters
                      between s/w failures
C
               iy - (input.) a vector of length nn, indicating the
C
                      type of error which is 1 if the i-th
                      failure is caused by an error due to
C
                      imperfect debussins, or 0 otherwith
C
                n - (input.) no. of observations for s/w failure time
C
               en - (output.) an estimate of parameter n
C
               ep - (output.) an estimate of parameter p
C
               er - (output.) an estimate of parameter lambda
C
             ecov - (output.) an estimate of variance-covariance
C
                      matrix (3x3)
C
C
      subroutine mle (t,iy,nn,en,ep,er,ecov)
      dimension t(nn), iy(nn), ecov(3,3)
                                        THIS PAGE IS BEST QUALITY PRACTICABLE
     ×1=0.0
     x2=0.0
      y=0.0
      do 5 i=1,nn
                                         TOO OURY THE ISHED TO DOG
      x1=x1+t(i)
      x2=x2+t(i)*float(i-1)
     w=wffloat(iy(i))
    5 continue
     en0=float(nn+1)
     ep0=1.0
     do 15 j=1,20
     x3=x1*en0-ep0*x2
     ×4=0.0
     x5=0.0
     do 25 i=1,nn
     x4=x4+float(1-iy(i))/(en0-float(i-1))**2
      x5=x5+float(1-iy(i))/(en0-float(i-1))
   25 continue
      f0=x5-x1*float(nn)/x3
      fn=-x4+x1*x1*float(nn)/x3/x3
      fp=-x1*x2*float(nn)/x3/x3
      ph0=float(nn)*(1.0-ep0)*x2/x3-y
      Phn=-(1.0-ep0)*float(nn)*x2*x1/x3/x3
      php=float(nn)*x2*(-1.0+(1.0-ep0)*x2/x3)/x3
      hk=fn*php-phn*fp
     hh=-(f0*php-ph0*fp)/hk
      xk=-(fn*ph0-phn*f0)/hk
      en0=en0+hh
      ep0=ep0+xk
```

TABLE C.9 (Continued)

```
print 100, j, en0, ep0
100 format(i5,2e15.5)
    if (amax1(abs(hh/en0),abs(xk/ep0)) .le. 0.00001) so to 11
15 continue
 11 en=en0
   ep=ep0
   er=float(nn)/x3
    rnn=0.0
    rpr=0.0
    rnr=0.0
   do 35 i=1,nn
    rnn=rnn+1.0/(en-float(i-1))/(en-ep*float(i-1))
    rpr=rpr+float(i-1)/(en-ep*float(i-1))
    rnr=rnr+1.0/(en-ep*float(i-1))
 35 continue
    rnr=rnr/er
    rpp=rpr/(1.0-ep)
    rpr=-rpr/er
    rrr=float(nn)/er/er
    TX=TNN*TPP*TTT-TNT*TPP*TNT-TNN*TPT*TPT
    ecov(1,1)=(rpp*rrr-rpr*rpr)/rx
    ecov(1,2)=rpr*rnr/rx
    ecov(1,3)=-rpp*rnr/rx
    ecov(2,2)=(rnn*rrr-rnr*rnr)/rx
    ecov(2,3)=-rnn*rpr/rx
    ecov(3,3)=rnn*rpp/rx
    ecov(2,1)=ecov(1,2)
    ecov(3,1)=ecov(1,3)
                                   THIS PAGE IS BEST QUALITY PRACTICABLE
    ecov(3,2)=ecov(2,3)
    return
                                   FROM COPY PARMISHED TO DDQ
   end
```

THIS PAGE IS BEST QUALITY PRACTICABLE

TABLE C.10 SUBROUTINE BAYES

```
subroutine bases (trisrnn,alpha,beta,pi,rho,xmu,samma,bn,bp,br)----
                  - obtain bayesian estimates of unknown parameters no po
  function
C
                       and lambda for idm
C
                  - call bases (tris,noralpha,beta,pi,rho,xmu,samma,bo,bp,br)
C
  usase
                t - (input.) a vector of length nn, containing time
   parameters
                       between s/w failures
               iy - (input.) a vector of length nn, indicating the
                       type of error which is 1 if the i-th
                       failure is caused by an error due to
                       imperfect debussins, or 0 otherwith
C
               nn - (input.) no. of observations for s/w failure time
C
            alpha - (input.) shape parameter of samma prior for n
             beta - (input.) scale parameter of samma prior for n
C
               pi - (input.) first parameter of beta prior for p
              rho - (input.) second parameter of beta prior for p
C
            xmu - (input.) shape parameter of damma prior for lambda
gamma - (input.) scale parameter of damma prior for lambda
C
C
               bn - (output.) bayesian estimate of n
C
               bp - (output.) bayesian estimate of p
C
C
               br - (output.) bayesian estimate of lambda
C
      subroutine bayes (t,iy,nn,alpha,beta,pi,rho,xmu,samma,bn,bp,br)
      dimension t(nn), iy(nn)
      x1=0.0
      x2=0.0
      ¥=0.0
      do 5 i=1,nn
      x1=x1+t(i)
      x2=x2+t(i)*float(i-1)
      y=y+float(iy(i))
    5 continue
      bn0=float(nn+1)
      bp0=1.0
      do 15 j≈1,20
      x3=x1*bn0-be0*x2+samma
      ×4=0.0
      x5=0.0
```

TABLE C.10 (Continued)

do 25 i=1,nn x4=x4+float(1-iy(i))/(bn0-float(i-1))**2 x5=x5+float(1-iy(i))/(bn0-float(i-1))25 continue f0=x5-x1*(xmu+float(nn-1))/x3+(alpha-1.0)/bn0-beta fn=-x4+x1*x1*(xmu+float(nn-1))/x3/x3-(alpha-1.0)/bn0/bn0 fp=-x4*x2*(xmu+float(nn-1))/x3/x3 PhO=(xmu+float(nn-1))*(1.0-bpO)*x2/x3-y+(pi-1.0)*(1.0-bpO)/bpO-(rho-1.0) Phn=-(1.0-bp0)*(xmu+float(nn-1))*x2*x1/x3/x3 Php=(xmutfloat(nn-1))*x2*(-1.0+(1.0-bp0)*x2/x3)/x3-(pi-1.0)/bp0/bp0 hk=fn*php-phn*fp hh=-(f0*php-ph0*fp)/hk xk=-(fn*ph0-phn*f0)/hk bn0=bn0+hh bro=bro+xk print 100, j, bn0, bp0 100 format(i5,2e15.5) if (amax1(abs(hh/bn0),abs(xk/bp0)) .le. 0.00001) so to 11 15 continue 11 bn=bn0 be=be0 br=(xmu+float(nn-1))/x3 return end

> THIS PAGE IS BEST QUALITY PRACTICABLE FROM OQPY FURMISHED TO DDQ

TABLE C.11 SUBROUTINE GGUB

```
subroutine saub (iseed,n,r)-----from imsl-----
  function - basic uniform (0.1) pseudo-random number
                     senerator
                  - call ssub (iseed,n,r)
  parameters iseed - (input.) an integer value in the exclusive
                      range (1,2147483647), iseed is replaced by
C
                       a new iseed to be used in subsequent calls.
                n - (input.) no. of deviates to be senerated
              r(n) - (output vector of length n, containing the
C
                      deviates in (0,1)
     subroutine ssub(iseed,n,r)
     dimension r(1)
     double precision z,dpm,dp
     dpm=dble(float(2**31-1))
     dp=dble(float(1/2**31))
     z=iseed
     do 5 i=1.n
     z=dmod(16807.d0*z,dpm)
   5 r(i)=z/dpm
     iseed=z
     return
     end
```

C.3 Programs for the Imperfect Maintenance Model of Section 3

Computer programs to compute the following quantities of interest for given N, p, $\,\lambda\,$ and $\,\mu\,$ are given in Tables C.12 to C.18.

- Mean and variance of the first passage time from N to n_0 (SUBROUTINE COMP)
- · Pdf and cdf of the first passage time from N to $^{\rm n}_{\rm 0}$ (SUBROUTINE FIRST)
- Probability of the software system being operational with n_0 remaining errors at time t (SUBROUTINE STT)
- Software system availability at time t (SUBROUTINE AVAIL)
- Expected number of errors detected and corrected by time t
 (SUBROUTINE EXPCT)

The programs are self-explanatory and include the required subroutines (MDGAMM and GAMMA).

TABLE C.12 SUBROUTINE COMP

```
subroutine comp (n,p,r,xm,k1,k2,et,vt,a,b,r1,r2)--
                          - compute mean and variance of first passage
   function
                              time, estimate the samma parameters, and
                              obtain the constants r1 and r2
.
                          - call comp (n:p:r:xm:k1:k2:et:vt:a:b:r1:r2)
C
                       n - (input.) initial no. of errors
C
   parameters
                         - (input.) prob. of perfect debussins
C
                        r - (input.) detection rate / error
C
                      xm - (input.) correction rate / error
k1 - (input.) first destination
C
C
                      k2 - (input.) second destination
c
                      et - (output.) mean first passage time
C
                       vt - (output.) variance of first passage time
C
                        a - (output.) shape parameter of samma distribution
C

    b - (output.) scale parameter of samma distribution
    r1 - (output.) smaller constant

C
C
                       r2 - (output.) larger constant
C
C
      subroutine comp (n.p.r.xm.k1.k2.et.vt.a.b.r1.r2)
      kk1=k1+1
      kk2=k2+1
      rr=sart((r+xm)**2-4.0*p*r*xm)
      r1=(r+xm-rr)/2.0
      r2=(r+xm+rr)/2.0
      er1=0.0
      er2=0.0
      vr1=0.0
                                                 BELS PAGE IS BEST QUALITY PRACTICABLE
      Vr2=0.0
      if (kk1 .st. n) so to 11
      do 5 i=kk1.n
                                                  THE PAGE IS BUST QUALITY PRAY
      er1=er1+1.0/float(i)
      vr1=vr1+1.0/float(i)/float(i)
    5 continue
   11 er1=er1/r1
      Vr1=Vr1/r1/r1
      if (kk2 .st. n) so to 22
      do 15 i=kk2.n
      er2=er2+1.0/float(i)
      vr2=vr2+1.0/float(i)/float(i)
   15 continue
   22 er2=er2/r2
      vr2=vr2/r2/r2
      et=er1+er2
      vt=vr1+vr2
      if (vt .ea. 0.0) so to 33
      b=et/vt
      so to 99
   33 b=0.0
   99 a=et*b
      return
```

end .

TABLE C.13 SUBROUTINE FIRST

```
subroutine first (n.p.r.xm.k1.k2.t.pdf.cdf.r1.r2)-----
C
                       - compute pdf and cdf of first passage time
C
   function
                           and also obtain the constants r1 and r2
C
  nesae
                  - call first (n,p,r,xm,k1,k2,t,pdf,cdf,r1,r2)
C
                     n - (input.) initial no. of errors
C
                    p - (input.) prob. of perfect debussins
r - (input.) detection rate / error
C
C
                    xm - (input.) correction rate / error
C
                    k1 - (input.) first destination
C
                    k2 - (input.) second destination
C
                   pdf - (output.) pdf of first passage time
C
                 P cdf - (output.) cdf of first passage time
C
                    r1 - (output.) smaller constant
C
                    r2 - (output.) larger constant
   read. subroutines - comp, mdsamm, samma
C
C
C
      subroutine first (n,p,r,xm,k1,k2,t,pdf,cdf,r1,r2)
      if (min0(k1,k2) .1t. 0) so to 22
      call comp (n.p.r.xm.kl.k2.et.vt.a.b.r1.r2)
      if (b .ea. 0.0) so to 11
      x=b*t
      x1=(a-1.0)*alos(x)
      x2=alsamma(a)
      x3=x1-x-x2
      if (x3 .1t. -88.0) so to 33
      pdf=exp(x3)
      so to 44
   33 pdf=0.0
   44 call mdsamm (x,a,cdf)
                                       THIS PAGE IS BEST QUALITY PRACTICABLE
      return
                                       FROM OOPY PURMISHED TO DDC
   11 cdf=1.0
      Pdf=1.0
      return
   22 cdf=0.0
      pdf=0.0
      return
      end
```

TABLE C.14 SUBROUTINE STT

```
subroutine stt (n,p,r,xm,nO,t,prob)-----
                       - compute the probability of being nO s/w
   function
                           errors remaining at time t
C
                       - call stt (n,p,r,xm,n0,t,prob)
  usade
C
                     n - (input.) initial no. of errors
                     P - (input.) prob. of perfect debussins
C
                    r - (input.) detection rate / error xm - (input.) correction rate / error
C
C
                    n0 - (input.) specified no. of errors
C
                    t - (input.) time
                  prob - (output.) prob. of being nO errors remaining
                           at time t
                       - first
   read. subroutine
C
C
      subroutine stt (n,p,r,xm,n0,t,prob)
      call first (n,p,r,xm,n0,n0,t,pdf,cdf,r1,r2)
      Prob=cdf
      k=n0-1
      c1=(r-r2)/(r1-r2)
```

prob=cdf
k=n0-1
c1=(r-r2)/(r1-r2)
c2=(r1-r)/(r1-r2)
call first (n,p,r,xm,k,n0,t,pdf,cdf,r1,r2)
prob=prob-c1*cdf
call first (n,p,r,xm,n0,k,t,pdf,cdf,r1,r2)
prob=prob-c2*cdf
return
end

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY PARELSHED TO DDG

TABLE C.15 SUBROUTINE AVAIL

```
subroutine avail (n.p.r.xm.t.at)-----
                      - compute s/w system availability at time t - call avail (n.p.,r.xm,t.at)
 function
C
  usase
C
  parameters
                    n - (input.) initial no. of errors
C
                    P - (input.) prob. of perfect debussins
                  r - (input.) detection rate / error
C
                   xm - (input.) correction rate / error
C
C
                   t - (input.) time
                   at - (output.) availability at time t
  read. subroutine - stt
C
C
     subroutine avail (n,p,r,xm,t,at)
     at=0.0
     n1=n+1
     do 5 i=1.n1
     n0=i-1
     call stt (n,p,r,xm,n0,t,prob)
      at=at+prob
    5 continue
      return
      end
```

THIS PAGE IS BEST QUALITY PRACTICABLE

TABLE C.16 SUBROUTINE EXPCT

```
subroutine exect (n,p,r,xm,t,xmd,xmc)-----
C
   function
                      - compute mean no. of errors detected and
C
                          corrected by time t
                    - call expct (n,p,r,xm,t,xmd,xmc)
  usade
C
                 n - (input.) initial no. of errors
  parameters
C
                   > - (input.) Prob. of Perfect debussins
                    r - (input.) detection rate / error
C
                   xm - (input.) correction rate / error
C
                    t - (input.) time
C
                  xmd - (output.) mean no. of errors detected by
c
C
                          time t
C
                  xmc - (output.) mean no. of errors corrected by
                          time t
C
C
   read. subroutine
                      - first
C
      subroutine exect (n,p,r,xm,t,xmd,xmc)
      h1=0.0
     h2=0.0
     h=0.0
      do 5 i=1,n
      k=i-1
      call first (n,p,r,xm,k,i,t,pdf,cdf,r1,r2)
      h1=h1+cdf
      call first (n,p,r,xm,i,k,t,pdf,cdf,r1,r2)
     h2=n2+cdf
      call first (n,p,r,xm,k,k,t,pdf,cdf,r1,r2)
      h=h+cdf
    5 continue
      xmc=h/p
xmd=(h1*(1.0-xm/r1)+h2*(xm/r2-1.0))*r/(r1-r2)
      return
      end
```

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY FURNISHED TO DDC

TABLE C.17 SUBROUTINE MDGAMM

```
subroutine mdsamm (x,p,prob)-----from ims1------
C
C
 function - compute imcomplete samma distribution
C
  usade
                     - call mdsamm (x,p,prob)
  parameters
                x - (input.) value to which samma is to be integrated
C
                  P - (input.) samma parameter
C
                Prob - (output.) prob.=integral of gamma(p) to x
C
  read. subroutines - samma
C
C
     subroutine mdsamm (x,p,prob)
     dimension v(6),v1(6)
     equivalence (v(3),v1(1))
     Prob=0.0
     if (x .se. 0.0) so to 5
     so to 9000
   5 if (p .st. 0.0) so to 10
     so to 9000
   10 if (x .ea. 0.0) so to 9005
     fnls=alsamma(p)
     cnt=p*alog(x)
     went=x+fnls
     if ((cnt-yent) .st. -88.0) so to 15
     ax=0.0
     so to 20
  15 ax=exp(cnt-ycnt)
  20 bis=1.e35
     cut=1.e-8
     cut=1.e-8
if ((x .le. 1.0) .or. (x .lt. p)) so to 40
     y=1.0-p
     z=x+4+1.0
     cnt=0.0
     v(1)=1.0
     v(2)=x
     v(3)=x+1.0
     v(4)=z*x
     Prob=v(3)/v(4)
   25 cnt=cnt+1.0
     y=y+1.0
     z=z+2.0
     went=w*ent
     v(5)=v1(1)*z~v(1)*sent
     v(6)=v1(2)*z-v(2)*yent
     if (v(6) .ea. 0.0) so to 50
```

STATE OF THE STATE OF THE PERSON STATES

TABLE C.17 (Continued)

ratio=v(5)/v(6) reduc=abs(prob-ratio) if (reduc .st. cut) so to 30 if (reduc .le. ratio*cut) so to 35 30 prob=ratio so to 50 35 prob=1.0-prob*ax so to 9005 40 ratio=p cnt=1.0 Prob=1.0 45 ratio=ratio+1.0 cnt=cnt*x/ratio Prob=Probtent if (cnt .st. cut) so to 45 Prob=Prob*ax/P so to 9005 50 do 55 i=1,4 v(i)=v1(i) 55 continue if (abs(v(5)) .1t. bis) so to 25 do 60 i=1,4 v(i)=v(i)/bis 60 continue so to 25 9000 continue 9005 return end

function alsamma(P)
if (P.st. 31.0) so to 15
call samma (P,SP)
alsamma=alos(SP)
return
15 z1=(P-0.5)*alos(P)-P+0.5*alos(2.0*3.1415)
z2=1.0/12.0/P
z3=1.0/360.0/P/P/P
z4=1.0/1260.0/P/P/P/P/P
z5=1.0/1680.0/P/P/P/P/P/P
alsamma=z1+z2-z3+z4-z5
return
end

THIS PAGE IS BEST QUALITY PRACTICABLE
THOS OUTY PAREISHED TO DDG

THIS PAGE IS BEST QUALITY PRACTICABLE PROM COPY PURHISHED TO DDC

TABLE C.18 SUBROUTINE GAMMA

```
subroutine samma (xx, dx) -----
                                                   -----from imsl-----
C
e function
                       - compute a samma function of parameter xx
  usade
                        - call samma (xx,sx)
C
                    xx - (input.) parameter of samma function sx - (output.) value of samma function
   parameters
C
C
C
.
      subroutine damma(xx,dx)
      1f(xx-57.) 6,6,4
    4 sx=1.0e30
      return
    6 x=xx
      err=1.e-6
      #x=1.
      if(x-2.) 50,50,15
   10 if (x-2.) 110,110,15
   15 x=x-1.
      SX=SXXX
   so to 10
50 if (x-1.) 60,120,110
   60 if (x-err) 62,62,80
   62 y=float(int(x))-x
   if (abs(y)-err) 120,120,64
64 if (1.-y-err) 120,120,70
   70 if(x-1.) 80,80,110
   X/xE=xE 08
      x=x+1.
      so to 70
  110 y=x-1.
      #y=1.+y*(-0.5771017+y*(0.9858540+y*(-0.8764218+y*(0.8328212+y*(-0.5684729
\c+y*(0.2548205+y*(-0.0514993)))))))
      EXXESXE
  120 return
      end
```

C.4 Program for Bayesian Software Correction Limit Policies
(Section 5)

Computer programs to compute the optimum policies for model 1 and model 2 are given in Tables C.19 to C.23. These programs perform the following functions:

- For model 1, simulate error occurrence time and correction time and then compute Bayesian estimates of mean correction time, optimum correction limit time, and its minimum cost per unit time (SUBROUTINE MDL1)
- For model 2, simulate error occurrence time and correction time and then compute Bayesian estimate of mean correction time, optimum ocrrection limit time, and its minimum cost per unit times, and also provide the optimum sample size (SUBROUTINE MDL2)

The programs are self-explanatory and include the required subroutines (eg. DATA 1, GGUB and OPTMM).

TABLE C.19 SUBROUTINE MDL1

```
subroutine mdl1 (iseed,nn,r,xmu1,xmu2,c,alpha,beta,xn,yn,tt,ct)------
  function
                 - simulate error occurrence time and correction time
C
                         and then compute bayesian estimate of mean
                         correction time, optimum correction limit time,
.
                         and its minimum cost per unit time
C
           - call mdl1 (iseed,nn,r,xmu1,xmu2,c,alpha,beta,xn,un,tt,ct)
C
C
  parameters iseed - (input.) an integer value in the exclusive range
                        (1,2147483647). iseed is replaced by a new iseed
                        to be used in subsequent calls.
                nn - (input.) no. of observations
C
                 r - (input.) error occurrence rate
              xmu1 - (input.) mean correction time in phase 1
C
              xmu2 - (input.) mean correction time in phase 2
                 c - (input.) a vector of length 3
                       c(1) - cost per unit time of error correction
                                in pase 1
C
                        c(2) - cost per unit time of error correction
                                in pase 2
C
                       c(3) - sampling cost per sample size
             alpha - (input.) a vector of length 2
C
                       alpha(1) - shape parameter of inverted samma
                                     prior for mean correction time
C
                                     in pase 1
                       alpha(2) - shape parameter of inverted samma
C
          prior for error occurrence rate
beta - (input.) a vector of length 2
C
C
                       beta(1) - scale parameter of inverted samma
C
                                    prior for mean correction time
C
                                    in Pase 1
C
                       beta(2) - scale parameter of inverted samma
                                    prior for error occurrence rate
C
                xn - (output.) bayesian estimate of error occurence rate
                 wn - (output.) bayesian estimate of mean correction
C
C
                       time in pase 1
                 tt - (output.) optimum correction limit time
C
                 ct - (output.) minimum cost per unit time
C
  read. subroutines - data1, optmm
     subroutine mdl1 (iseed,nn,r,xmu1,xmu2,c,alpha,beta,xn,yn,tt,ct)
     dimension c(3),alpha(2),beta(2),xy(2)
     call data1 (iseed,nn,r,xmu1,xx,yy)
     xn=(xx+beta(2))/(alpha(2)+float(nn-1))
     a=0.0
     b=xn
     call optmm (nn, yy, c, alpha, beta, xmu2, a, b, yn, tt, ct)
     return
     end
```

AD-A067 690

SYRACUSE UNIV N Y

BAYESIAN SOFTWARE PREDICTION MODELS. VOLUME V. SUMMARY OF TECHN-ETC(U)

OCT 78 A L GOEL

F30602-76-C-0097 OCT 78 A L GOEL TR-78-5

UNCLASSIFIED

RADC-TR-78-155-VOL-5

NL

2 OF 2 ADA 067690







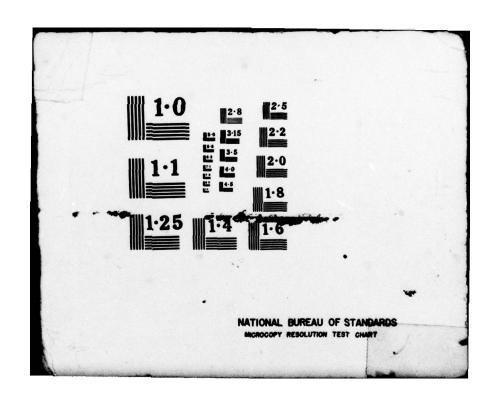








6 79



THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY PARELISHED TO BOQ

TABLE C.20 SUBROUTINE MDL2

```
subroutine md12 (iseed, r, xmu1, xmu2, c, alpha, beta, nn, xn, yn, tt, ct)--
  function - simulate error occurrence time and correction time
                         and then compute bayesian estimate of mean
                         correction time, optimum correction limit time,
                         and its minimum cost per unit time, and also
C
                         provide optimum sample size
C
                   - call mdl2 (iseed, r, xmu1, xmu2, c, alpha, beta, nn, xn, un, tt, ct)
  parameters iseed - (input.) an integer value in the exclusive range
c
                       (1,2147483647). iseed is replaced by a new iseed
                       to be used in subsequent calls.
C
                  r - (input.) error occurrence rate
              xmu1 - (input.) mean correction time in phase 1
C
              xmu2 - (input.) mean correction time in phase 2
C
                 c - (input.) a vector of length 3
.
                       c(1) - cost per unit time of error correction
                                in Pase 1
                       c(2) - cost per unit time of error correction
C
                                in Pase 2
C
                       c(3) - sampling cost per sample size
c
       alpha - (input.) a vector of length 2
C
                       alpha(1) - shape parameter of inverted samma
C
                                    prior for mean correction time
C
                                    in Pase 1
                       alpha(2) - shape parameter of inverted samma
                                    prior for error occurrence rate
C
        beta - (input.) a vector of length 2
.
                     beta(1) - scale parameter of inverted samma
C
                                   Prior for mean correction time
                                    in pase 1
C
                      beta(2) - scale parameter of inverted gamma
C
                                    prior for error occurrence rate
C
       nn - (output.) optimum sample size
                xn - (output.) bayesian estimate of ereor occurence rate
C
C
               wn - (output.) bayesian estimate of mean correction
                       time in pase 1
C
                tt - (output.) optimum correction limit time
                ct - (output.) minimum cost per unit time .
   read. subroutines - ssub, ortmm
```

subroutine mdl2 (iseed,r,xmu1,xmu2,c,alpha,beta,nn,xn,yn,tt,ct)

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY PAREISHED TO DDC

TABLE C.20 (Continued)

dimension c(3),alpha(2),beta(2),xy(2) XX=0.0 ww=0.0 do 5 nn=1,100 call ssub (iseed,2,xy) XX=XX-r*alog(XY(1)) YY=YY-xmu1*alog(xy(2)) if (nn .eq. 1) so to 5
xn=(xx+beta(2))/(alpha(2)+float(nn-1)) a=c(3)*float(nn) b=xxtxntaa call optmm (nn,yy,c,alpha,beta,xmu2,a,b,yn,t1,c1) a=c(3)*float(nn+1) b=xx+2.0*xn+yy+yn call optmm (nn,yy,c,alpha,beta,xmu2,a,b,yn,t2,c2) Print 100,nn,xn,yn,t1,c1,t2,c2 100 format(i5,6e15.5) if (t1-t2) 11,22,5 22 if (c1-c2) 11,11,5 5 continue 11 tt=t1 ct=c1 return end

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY FAMILISHED TO DDQ

TABLE C.21 SUBROUTINE DATAL

```
subroutine data1 (iseed,nn,r,xmu1,xx,yy)-----
   function
                       - simulate error occurrence time and correction
C
                           time in phase 1
                       - call data1 (iseed,nn,r,xmu1,xx,yy)
C
   usade
   parameters iseed - (input.) an integer value in the exclusive
C
                            ranse (1,2147483647). iseed is replaced
C
                           by a new iseed to be used in subsequent calls.
C
                    nn - (input.) sample size
C
                     r - (input.) error occurrence rate
C
                  xmu1 - (input.) mean correction time in phase 1
C
                    xx - (output.) total amount of error occurrence time vy - (output.) total amount of error correction time
C
C
                            in phase 1
   read. subroutine - saub
C
C
      subroutine data1 (iseed,nn,r,xmu1,xx,yy)
      dimension xy(2)
      xx=0.0
      WY=0.0
      do 5 i=1,nn
      call ssub (iseed,2,xy)
      xx=xx-r*alog(xy(1))
      yy=yy-xmu1*alos(xy(2))
    5 continue
      return
      end
```

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY PARKISHED TO DDC

DWYSO TELESCOPEROR TELESCOPEROR

TABLE C.22 SUBROUTINE GGUB

subroutine ssub (iseed,n,r)-----from imsl-----C function - basic uniform (0,1) pseudo-random number C senerator - call ssub (iseed,n,r) parameters iseed - (input.) an integer value in the exclusive ranse (1,2147483647), iseed is replaced by C a new iseed to be used in subsequent calls. C C n - (input.) no. of deviates to be senerated r - output vector of length n, containing the C deviates in (0,1) C C

subroutine ssub(iseed,n,r)
dimension r(1)
double precision z,dpm,dpn,dp
data dpm,dpn/2147483647.d0,1.d0/
dp=dpn/(dpm+dpn)
z=iseed
do 5 i=1,n
z=dmod(16807.d0*z,dpm,)
5 r(i)=z*dp
iseed=z
return
end

1

TOO Permate Covers. St.

CAND DESCRIPTION OF THE PROPERTY OF THE

TABLE C.23 SUBROUTINE OPTMM

```
subroutine optmm (nn:yy:c:alpha:beta:xmu2:a:b:yn:tt:ct)-----
                       - compute basesian estimate of mean correction
C
   function
                            time, optimum correction limit time, and its minimum cost per unit time
                       - call optmm (nn, yy, c, alpha, beta, xmu2, a, b, yn, tt, ct)
C
   usade
                    nn - (input.) no. of observations
ww - (input.) total amount of observed correction
   parameters
                           time
C
                     c - (input.) a vector of length 3
c
                           c(1) - cost per unit time of error correction
C
                                     in pase 1
C
                           c(2) - cost per unit time of error correction
C
                                      in pase 2
                           c(3) - sampling cost per sample size
C
                 alpha - (input.) a vector of length 2
                           alpha(1) - shape parameter of inverted samma
C
                                          prior for mean correction time
C
                                          in pase 1
                           alpha(2) - shape parameter of inverted samma
C
                                          prior for error occurrence rate
C
C
                  bets - (input.) a vector of length 2
                            beta(1) - scale parameter of inverted samma
C
                                         prior for mean correction time
C
C
                                         in Pase 1
                            beta(2) - scale parameter of inverted samma
C
C
                                         prior for error occurrence rate
                  xmu2 - (input.) mean correction time in pase 2
C
                     a - (input.) constant of a in equation (a-1)
C
                     b - (input.) constant of b in equation (a-1)
C
                    yn - (output.) bayesian estimate of mean correction
C
C
                            time in pase 1
                    tt - (output.) optimum correction limit time
C
C
                    ct - (output.) minimum cost per unit time
C
      subroutine ortmm (nn, yy, c, alpha, beta, xmu2, a, b, yn, tt, ct)
      dimension c(3),alpha(2),beta(2)
      aa=c(2)*b-a
      bb=(c(1)*b-a)/\times mu2
      #1=yytheta(1)
      =0=alpha(1)+float(nn-1)
      yn=41/20
      #2=bb/(aa+yn*c(2)-c(1))
      #3=(a0+1.0)/s1
      t=(s3/s2-1.0)*s1
      if (t .1t. 0.0) so to 22
      do 5 i=1,20
      rt=s3/(1.0+t/s1)
      ##t=#1/a0*(1.0-(1.0+t/#1)**(-a0))
      st=(1.0+t/s1)**(-a0-1.0)
      x1=aa+(c(2)-c(1))*ggt
      f0=rt*x1+(c(2)-c(1))*st-bb
      rrt=-s3/s1/(1.0+t/s1)**2
                                                THIS PAGE IS BEST QUALITY PRACTICABLE
      ff=rrt*x1
      h=-f0/ff
                                                PROS CORY PARKISHED TO DOG
      t=t+h
      Print 100,i,t
  100 format(15,e15.5)
```

TABLE C.23 (Continued)

if (abs(h/t) .le. 0.00001) so to 11
5 continue
11 tt=t
 ct=(c(1)-c(2)*xmu2*rt)/(1.0-xmu2*rt)
 return
22 tt=0.0
 ct=(atc(2)*xmu2)/(btxmu2)
 return
end

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY PUREISHED TO DDC and the second of the second